

DPST1091 / CPTG1391
Introduction to Programming
Week 9 – Lecture 2

Lecturer and Course Convener:

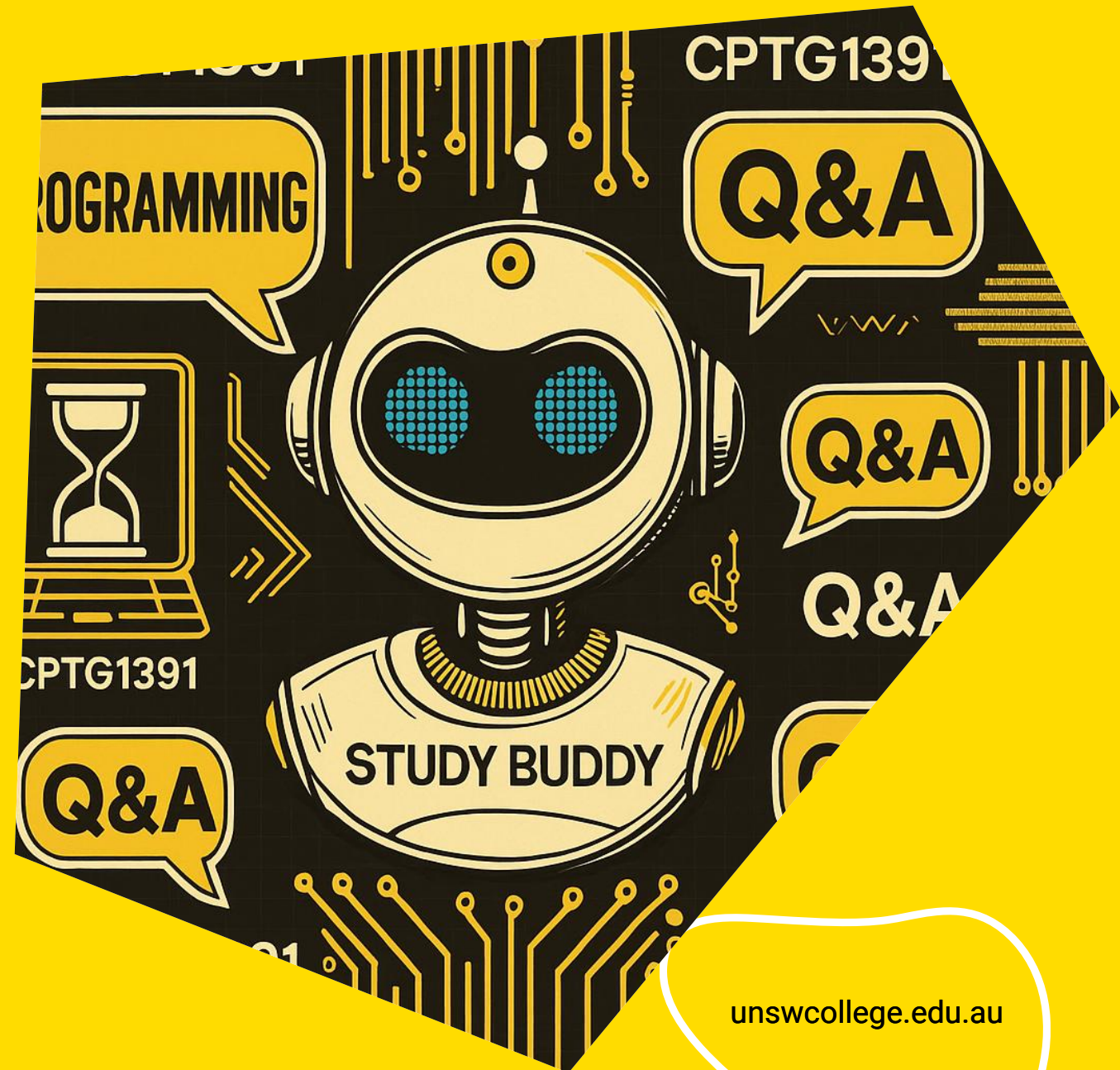
Dr Pantea Aria



UNSW
College

Linked Lists Functions

(part 1)




unswcollege.edu.au

Agenda

- **Last lecture**
Introduction to Linked Lists
- **Today**
 - Information about Practice Exam
 - Linked Lists Functions – Traverse and Insert

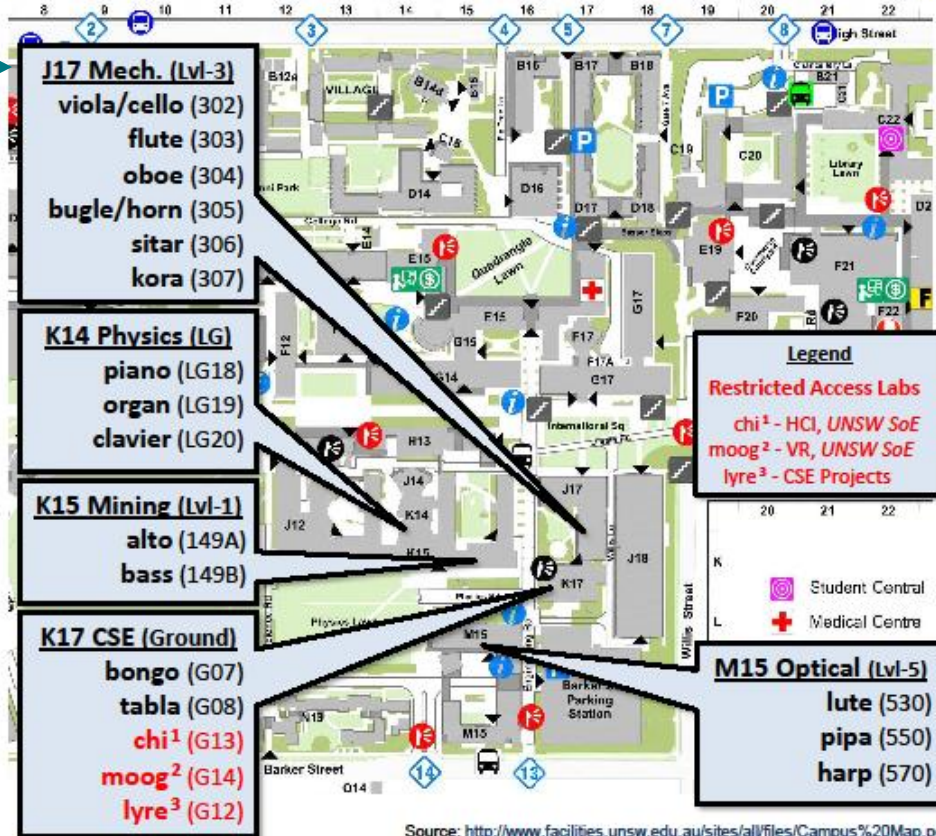
Practice Exam information

here →



School of Computer Science and Engineering
CSE Teaching Lab Locations

Kensington Campus



J17 Mech. (Lvl-3)
viola/cello (302)
flute (303)
oboe (304)
bugle/horn (305)
sitar (306)
kora (307)

K14 Physics (LG)
piano (LG18)
organ (LG19)
clavier (LG20)

K15 Mining (Lvl-1)
alto (149A)
bass (149B)

K17 CSE (Ground)
bongo (G07)
tabla (G08)
chi¹ (G13)
moog² (G14)
lyre³ (G12)

Legend


Restricted Access Labs
chi¹ - HCI, UNSW SoE
moog² - VR, UNSW SoE
lyre³ - CSE Projects

Student Central
Medical Centre

M15 Optical (Lvl-5)
lute (530)
pipa (550)
harp (570)

Source: <http://www.facilities.unsw.edu.au/sites/all/files/Campus%20Map.pdf>

CSE LAB ARE ADMINISTERED BY THE CSE 'COMPUTER SUPPORT GROUP'



There will be no lecture on Monday Week 10

It will be running during **next Monday lecture**

Held in **CSE Labs** for 2 hours (**Actual exam is about 3 hours**)

When: Monday 23rd March, 9:00 AM to 11:00 AM (be there 10 min early)

Where: J17 Ainsworth Building

Level3 <https://www.learningenvironments.unsw.edu.au/physical-spaces/k-j17>

Don't miss the chance to see what the exam environment is like and get used to using it. And try out some hurdle questions and other exam questions in an exam environment.

Don't worry, I won't mark it. It is only about you getting familiar with the exam environment.

Linked Lists, What are they?



An alternative way to store a collection of data is to use a **linked list**.

Arrays are extremely **useful** and we are definitely not replacing them.

A **linked list** is simply **another option** that may be more suitable in **certain situations**.

Linked lists work well for **sequential data**, such as:

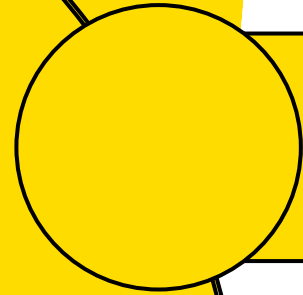
a to-do list

a train made up of connected carriages

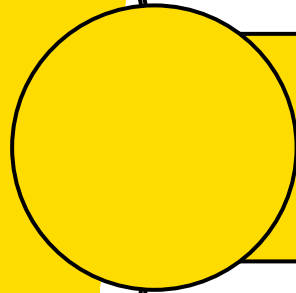
a queue of customers waiting in line

a chain of connected tasks in a workflow

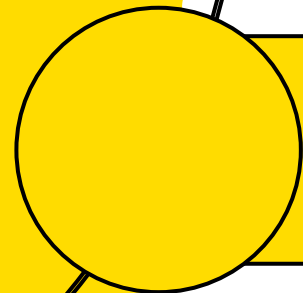
Arrays Advantages



Data elements are stored next to each other in a continuous (contiguous) section of memory.



They are efficient for both sequential access and direct (random) access.



Adding or removing elements at the end of the structure is simple and efficient.

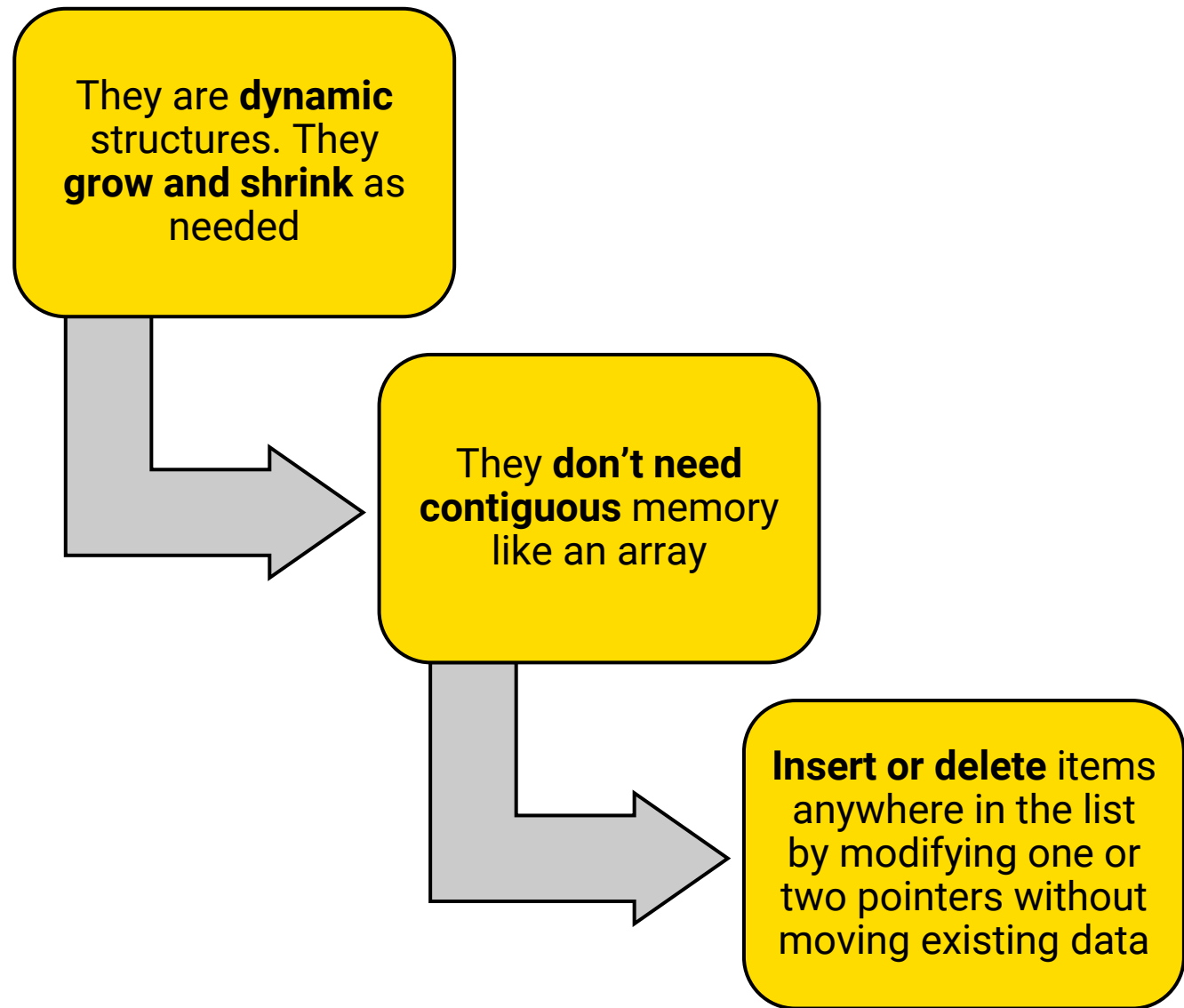
Arrays Disadvantages

Inserting or deleting elements in the middle can be inefficient and cumbersome.

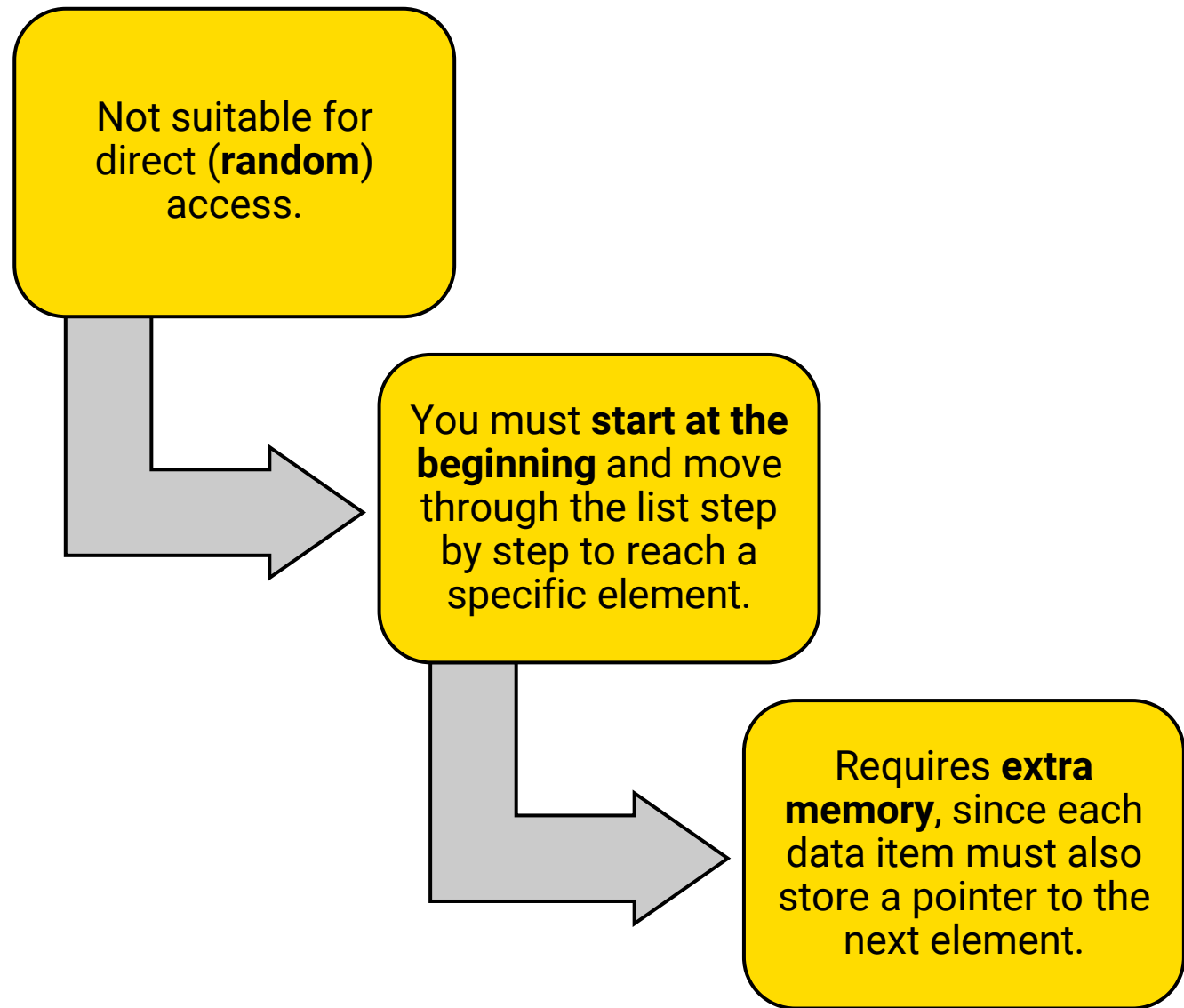
All the elements that come after the insertion or deletion point must be shifted to make space or close the gap.

If the array is static (fixed size), we actually cannot insert a new element once it is full.

Linked lists Advantages



Linked lists Disadvantages




Arrays and Linked Lists in the Memory

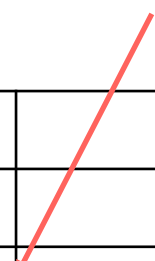
The elements in a linked list are not stored in contiguous memory locations.

Instead, they are distributed (scattered) across different areas of memory and connected using pointers

```
int numbers[] = {10, -4, 0, 12, 56};
```



0x30	
0x34	
0x38	10
0x3c	-4
0x40	0
0x44	12
0x48	56
0x4c	
0x50	
0x54	



0x30	
0x34	
0x38	10
0x3c	0x6c
0x40	
0x44	12
0x48	0x78
0x4c	
0x50	
0x54	

list

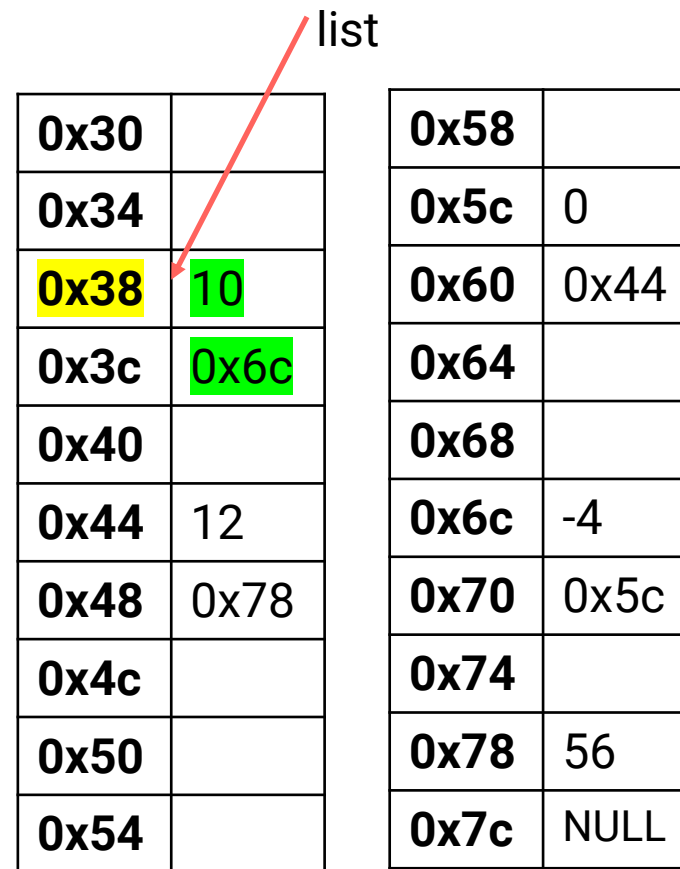
0x58	
0x5c	0
0x60	0x44
0x64	
0x68	
0x6c	-4
0x70	0x5c
0x74	
0x78	56
0x7c	NULL

Linked Lists in the Memory

You need a pointer to the **first node** in the list (called the **head**).

Each node stores its data along with a **pointer to the next node** in the sequence.

To access the full list, you follow each pointer **step by step**, like moving from one train carriage to the next in a connected train.

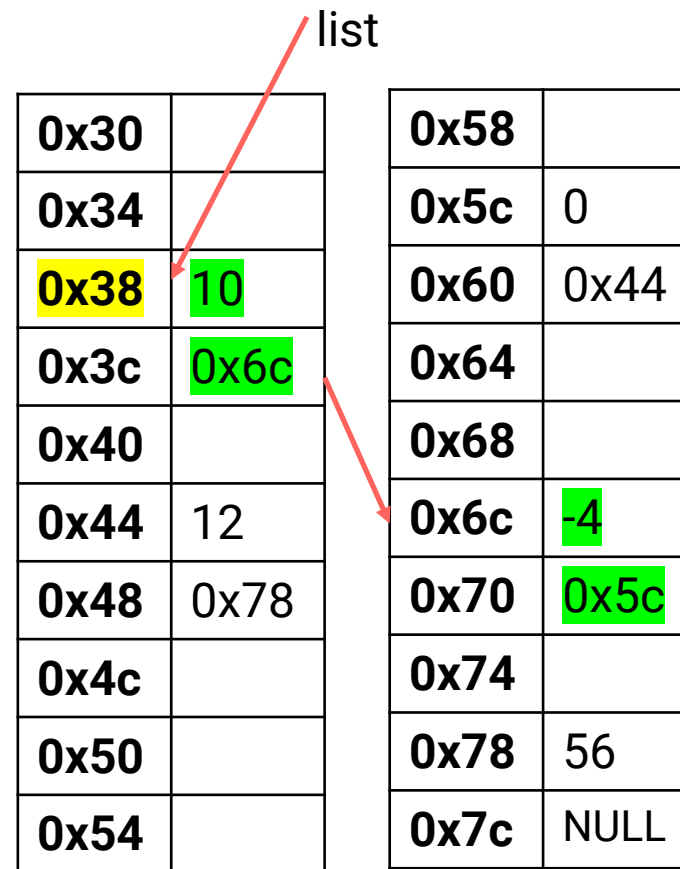


Linked Lists in the Memory

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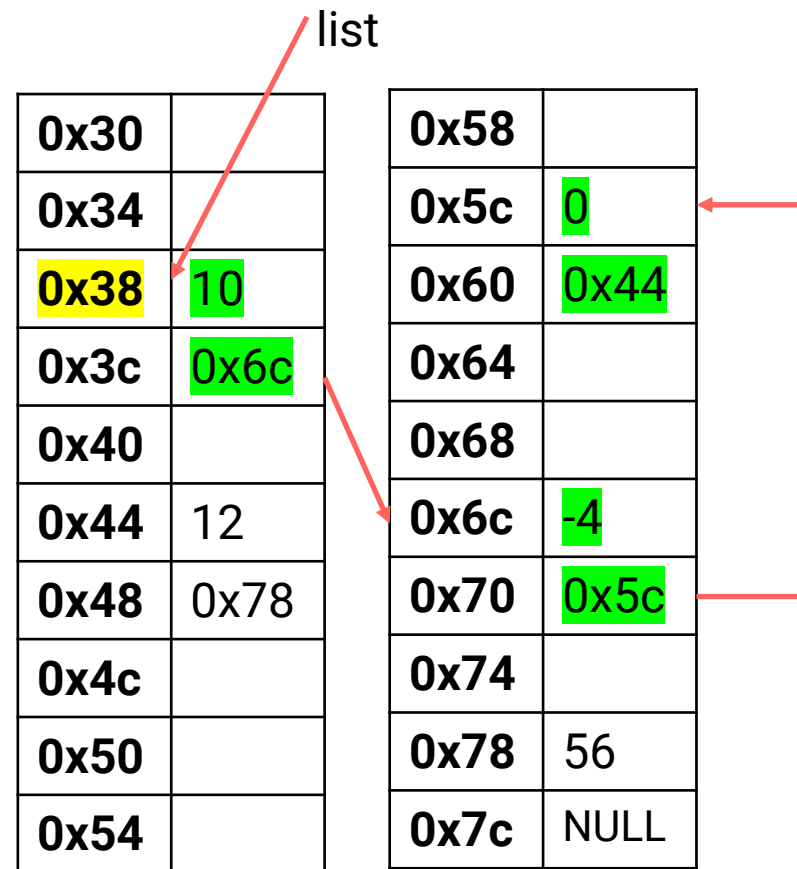


Linked Lists in the Memory

You need a pointer to the **first node** in the list (called the **head**).

Each **node** stores its **data** along with a **pointer** to the **next node** in the sequence.

To access the full list, you follow each pointer **step by step**, like moving from one train carriage to the next in a connected train.

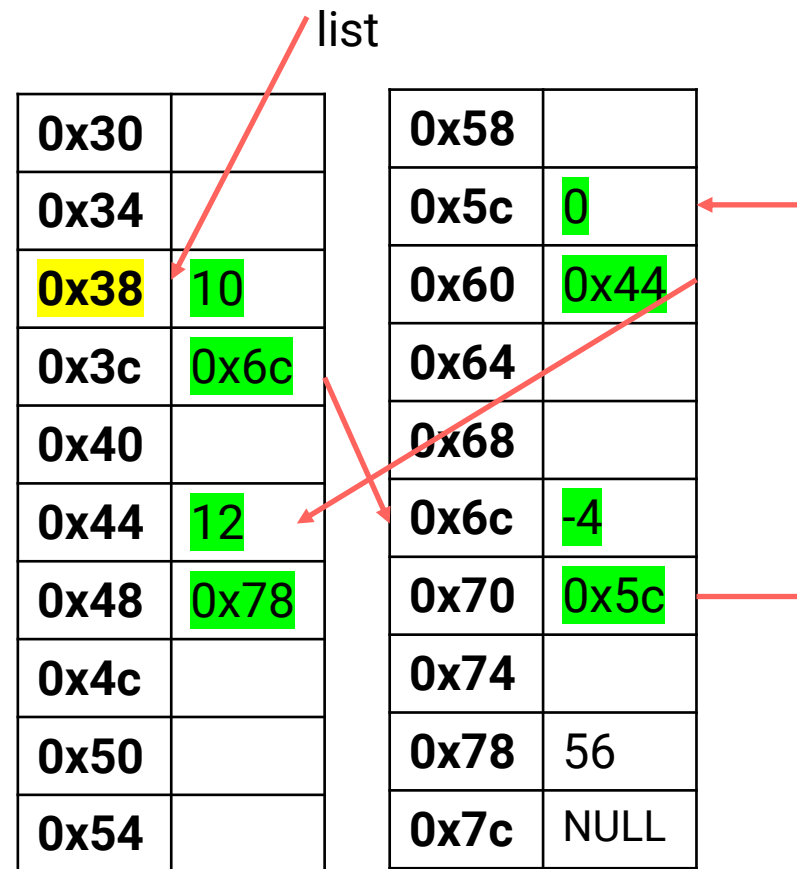


Linked Lists in the Memory

You need a pointer to the **first node** in the list (called the **head**).

Each node stores its data along with a **pointer to the next node** in the sequence.

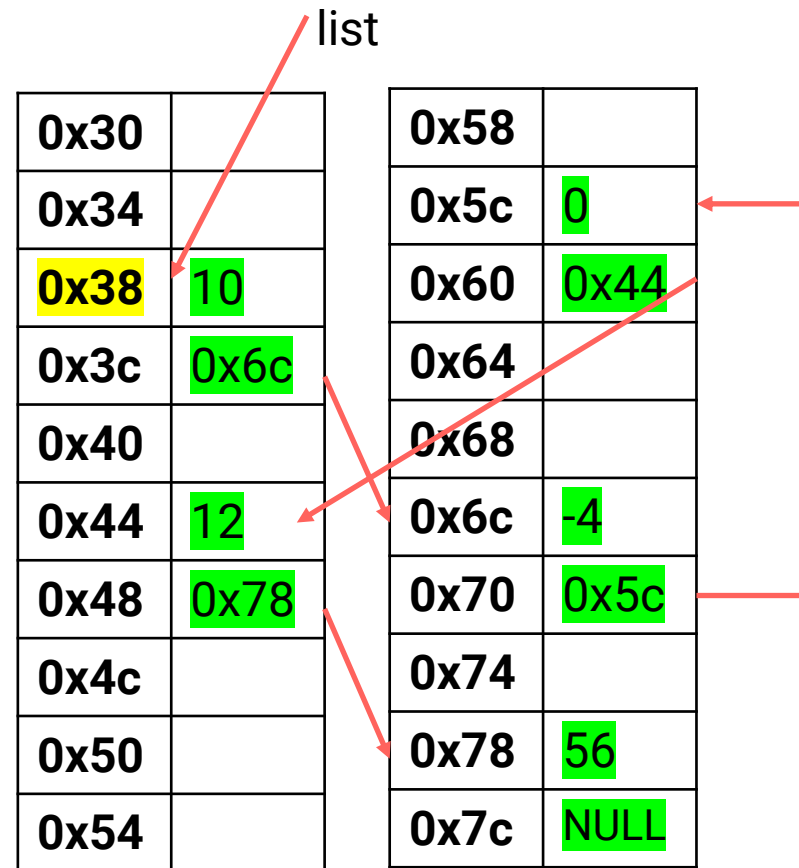
To access the full list, you follow each pointer **step by step**, like moving from one train carriage to the next in a connected train.



When the pointer to the next node contains **NULL**, it indicates that there are no more nodes to follow and you have reached the end of the list.

At this point, you have successfully traversed the entire linked list.

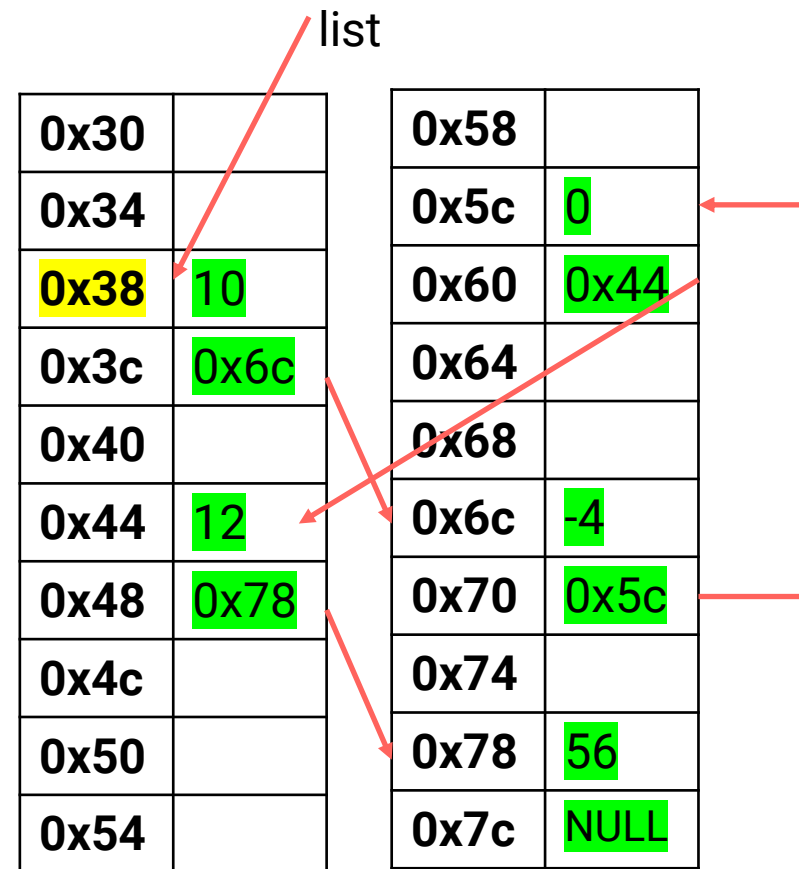
Linked Lists in the Memory



Linked Lists in the Memory

A linked list is described as sequential because we must begin at the first node and move through each node one by one to access elements.

Unlike arrays, we cannot directly access a specific element using an index; instead, we must traverse the list to reach it.

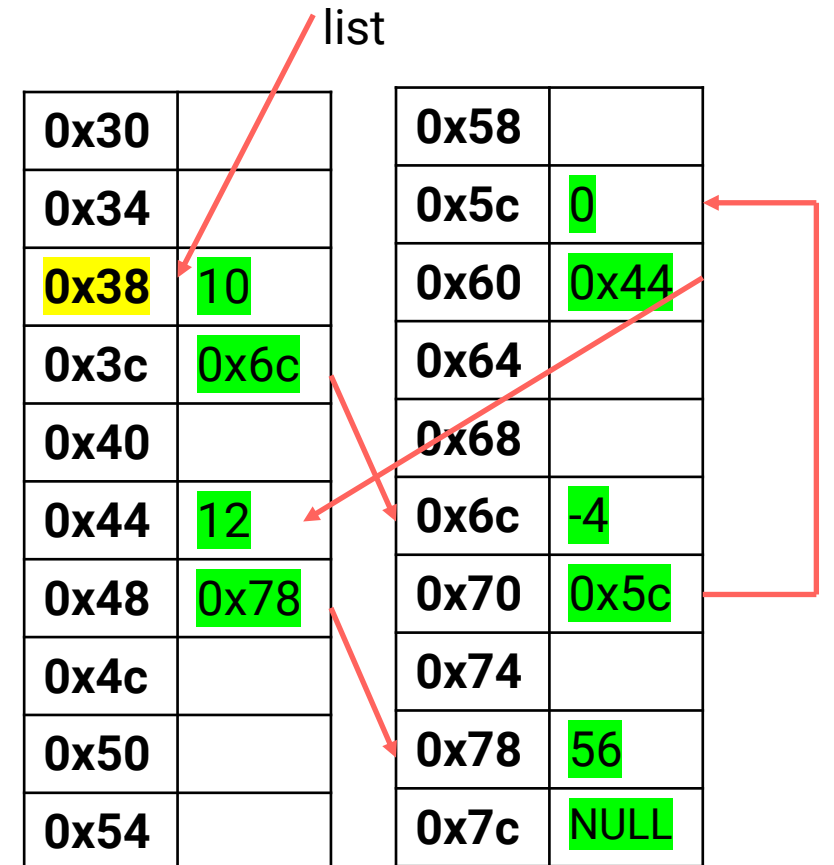


Nodes

In C, we can define a structure (**struct**) that allows us to **store both**:

- an **int value** (the data), and
- a **pointer containing the address of the next node** in the list.

```
struct node {  
    int data;  
    struct node *next;  
};
```



The variable **list** stores the address of the first node in the linked list

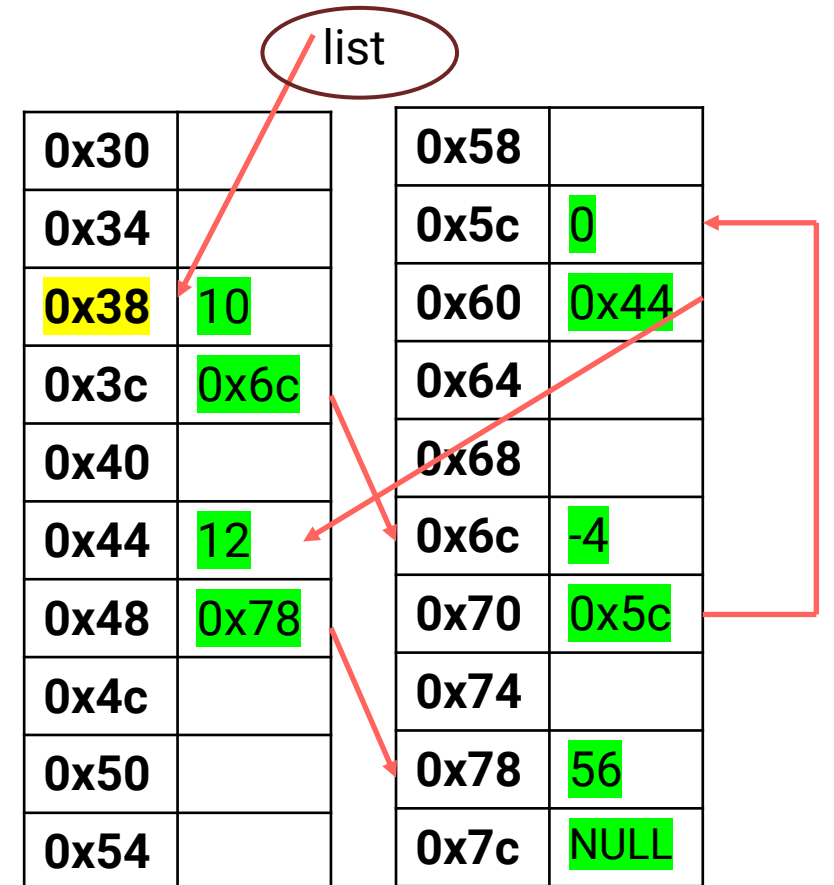
```

struct node *list;

struct node {
    int data;

    struct node *next;
};

```



Each node contains a **data** field.

```
struct node *list;

struct node {
    int data;

    struct node *next;
};
```

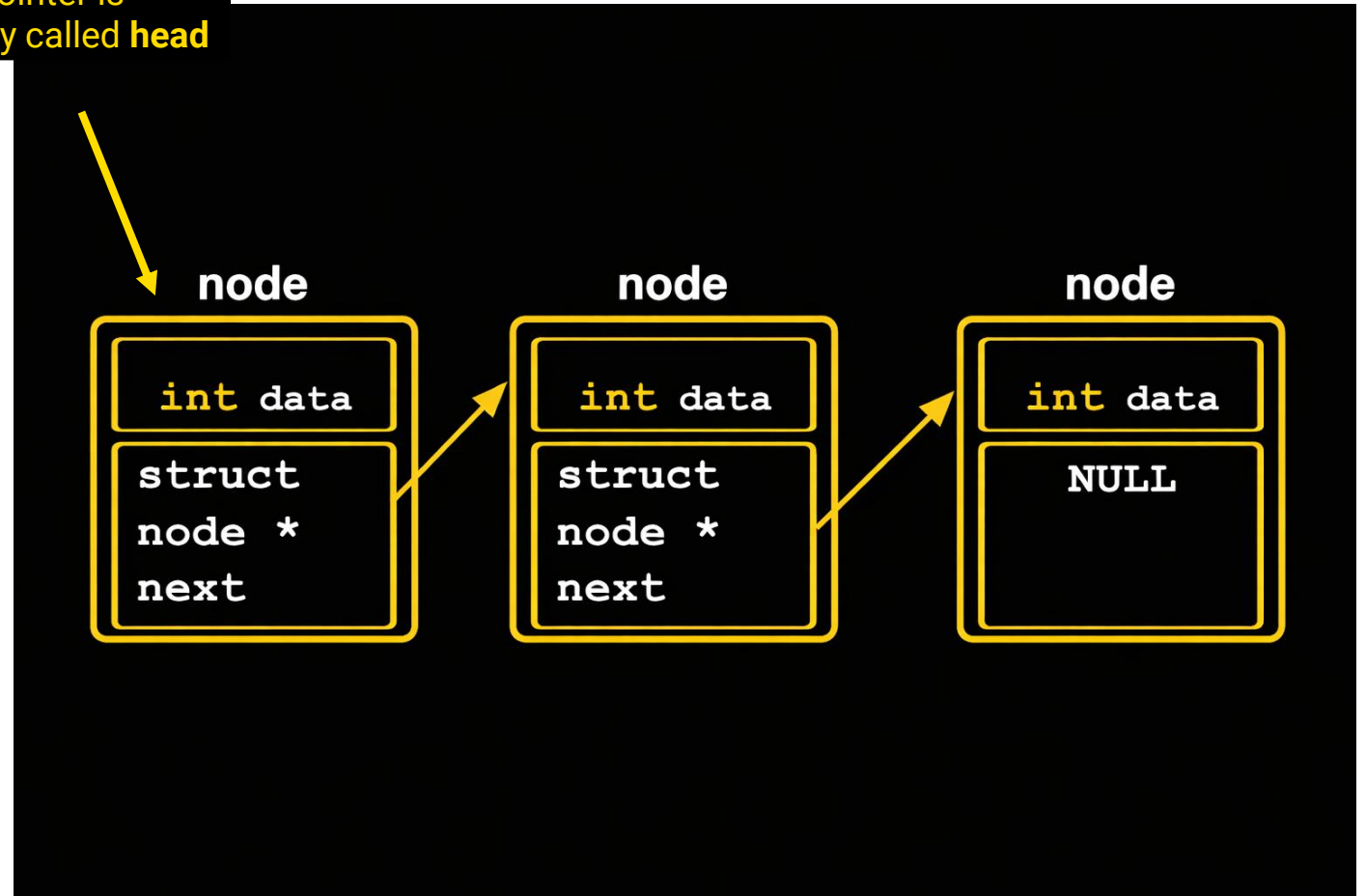
In this example, the **data** is a single **int**, but it can be **any data type** depending on your needs.

Later, we will explore linked lists that store **different types** of data.

Each node also includes a **pointer to the next node** in the list (which is of the **same type**).

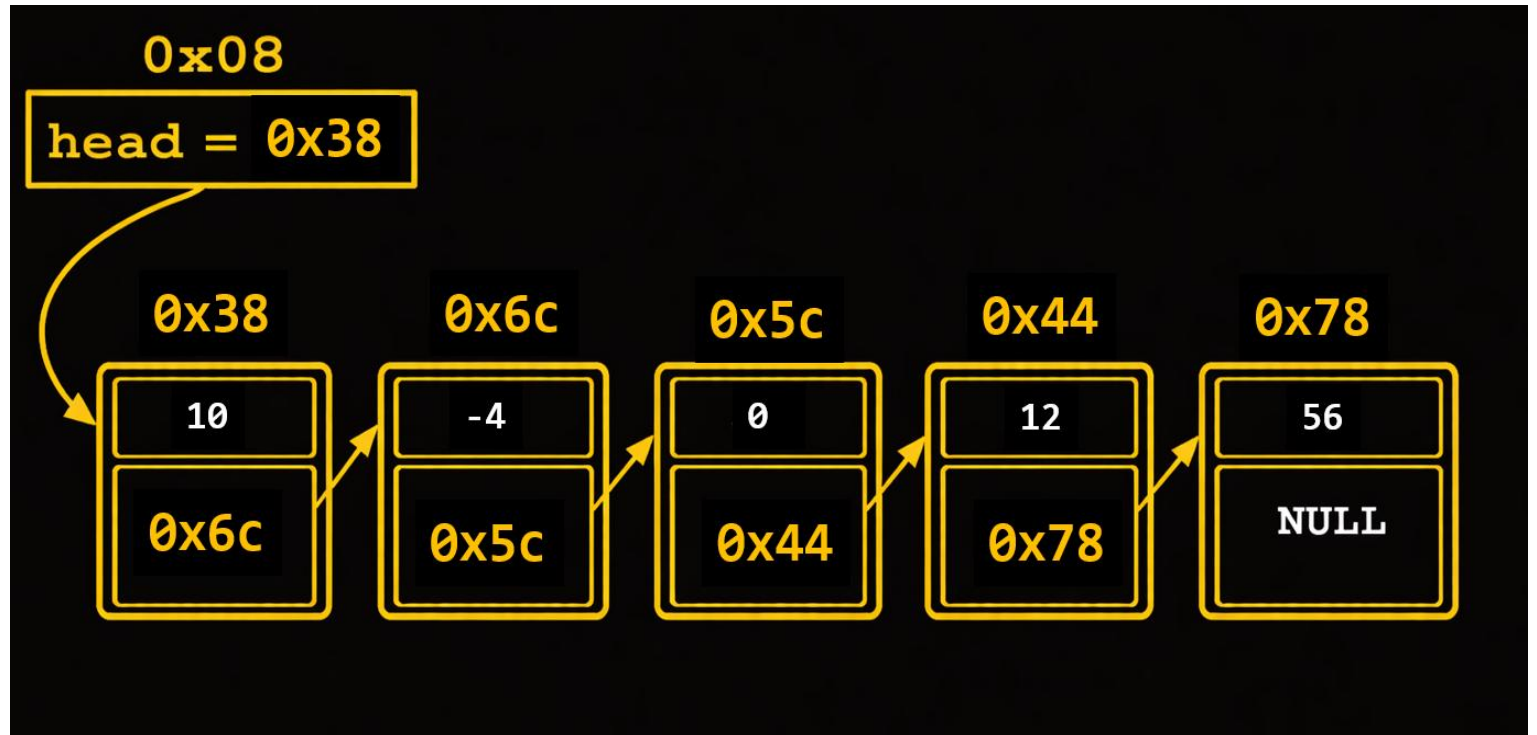
Visualising a Linked List

A pointer to the first node in the linked list (this pointer is commonly called **head**)



0x08
head = 0x38

0x30		0x58	
0x34		0x5c	0
0x38	10	0x60	0x44
0x3c	0x6c	0x64	
0x40		0x68	
0x44	12	0x6c	-4
0x48	0x78	0x70	0x5c
0x4c		0x74	
0x50		0x78	56
0x54		0x7c	NULL



Creating a Linked List

Here is the code to create a linked list with nothing in it.

```
struct node *head = NULL;
```

0x08

Head = NULL

Creating a Node

We use **malloc** to **allocate new nodes** on the **heap**.

This gives us full control over **memory management**, allowing us to:

create new nodes whenever needed, and

free them when they are **no longer required**.

Steps to Create a Node:

Allocate memory for a struct node using **malloc**.

Assign a value to the node's **data** field.

Set the **next pointer** to point to the **appropriate** node (or **NULL**).

Creating a List with 1 Node

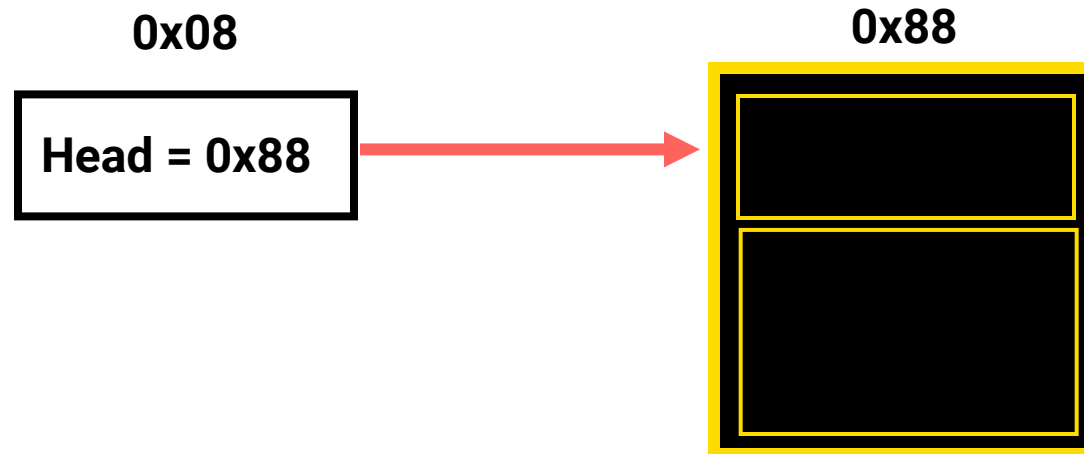
```
struct node {  
    int data;  
    struct node *next;  
};  
  
struct node *head = NULL;
```

0x08

Head = NULL

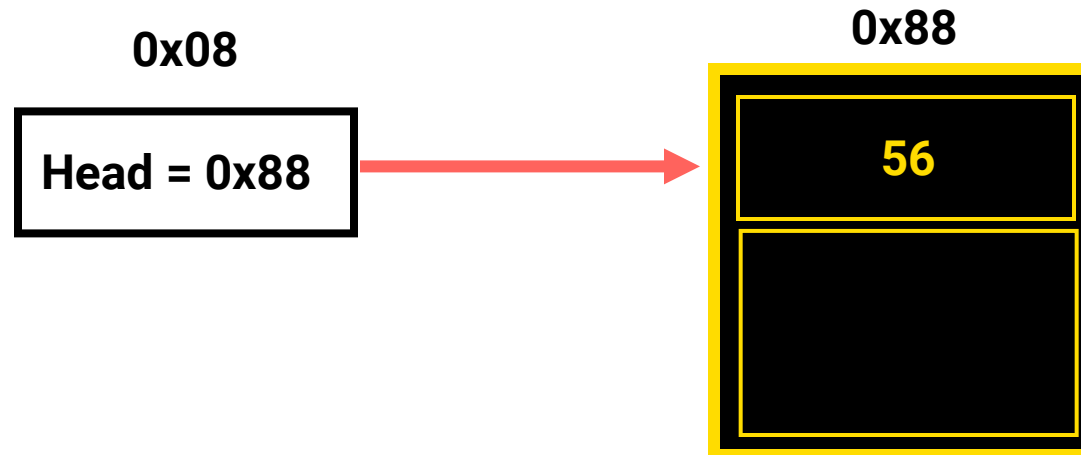
Creating a List with 1 Node

```
struct node {  
    int data;  
    struct node *next;  
};  
  
struct node *head = NULL;  
head = malloc(sizeof(struct node));
```



Creating a List with 1 Node

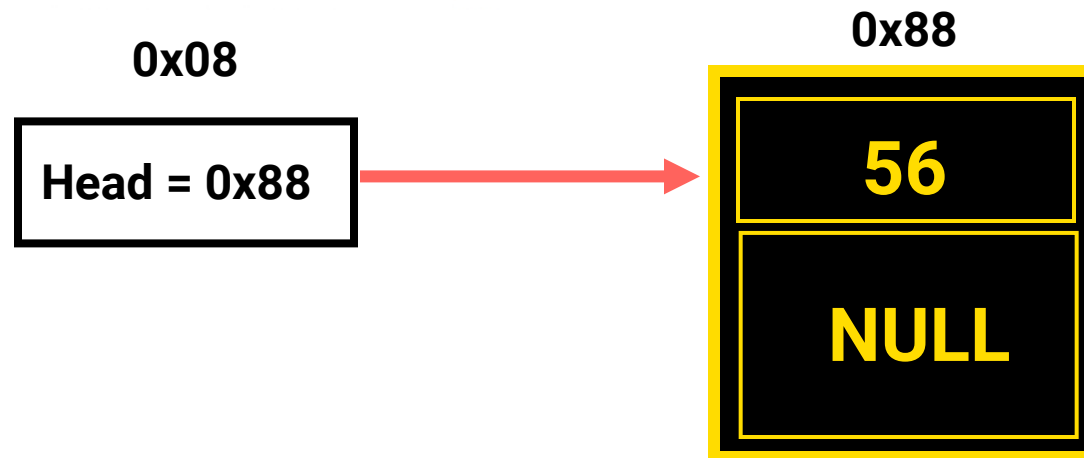
```
struct node {  
    int data;  
    struct node *next;  
};  
  
struct node *head = NULL;  
head = malloc(sizeof(struct node));  
head->data = 56;
```



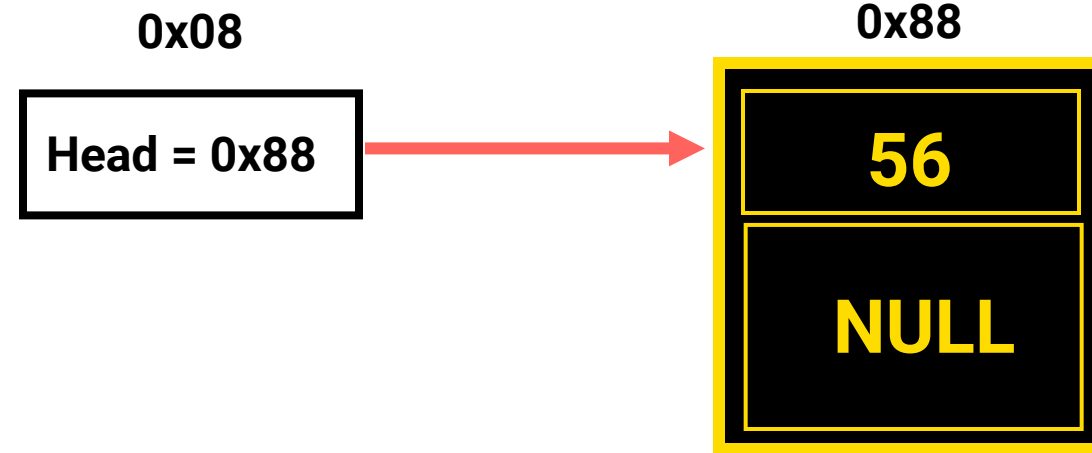
Creating a List with 1 Node

This is a linked list of size 1.

```
struct node {  
    int data;  
    struct node *next;  
};  
  
struct node *head = NULL;  
head = malloc(sizeof(struct node));  
head->data = 56;  
head->next = NULL;
```



Adding a new node to the end of our list

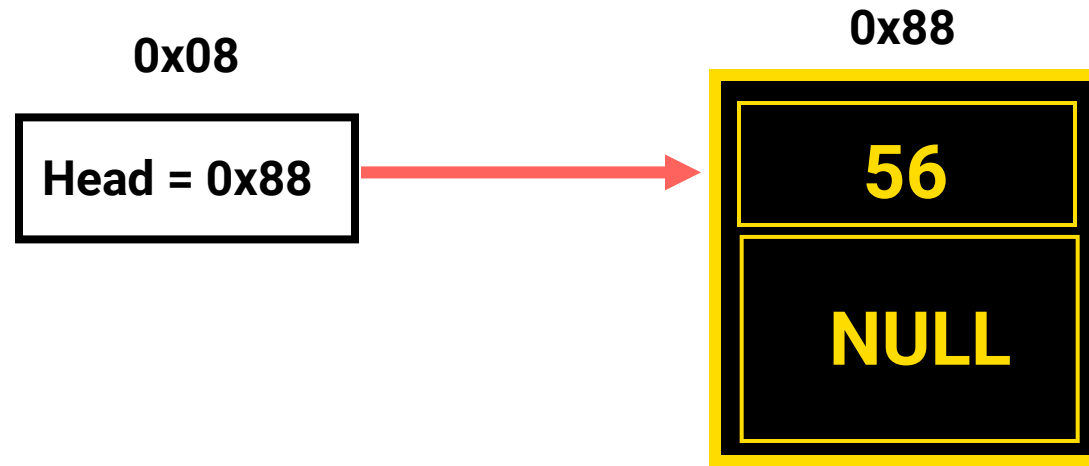


We will create a new node and attach it to the **end of the existing list.**

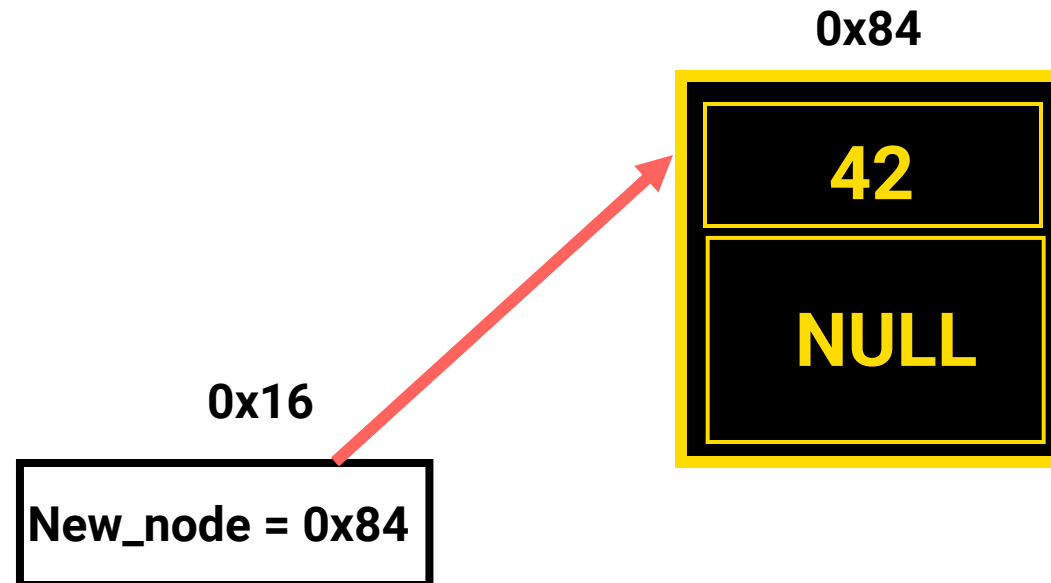
The **last node** in a linked list is commonly referred to as the **tail.**

Adding a new node to the end of our list

is the current node connected to the new node? (not yet)

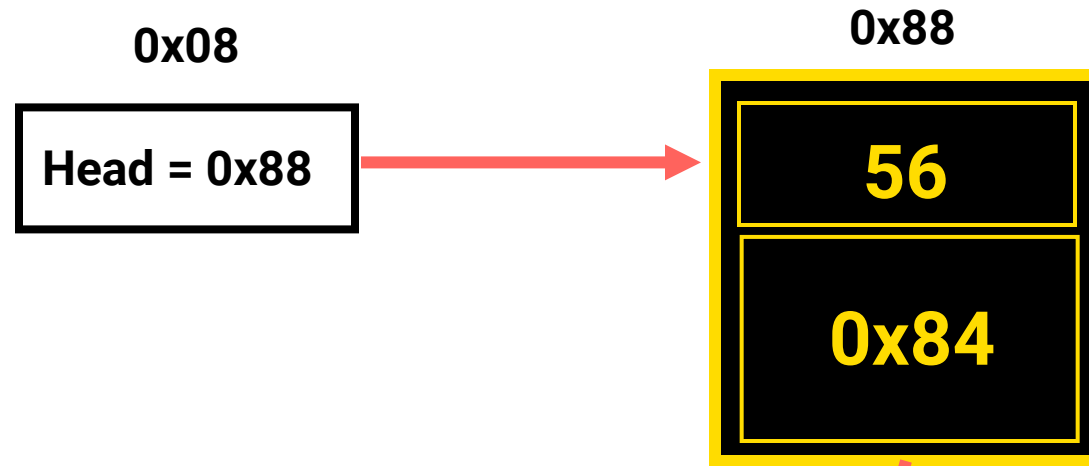


```
struct node *new_node = malloc(sizeof(struct node));  
new_node->data = 42;  
new_node->next = NULL;
```

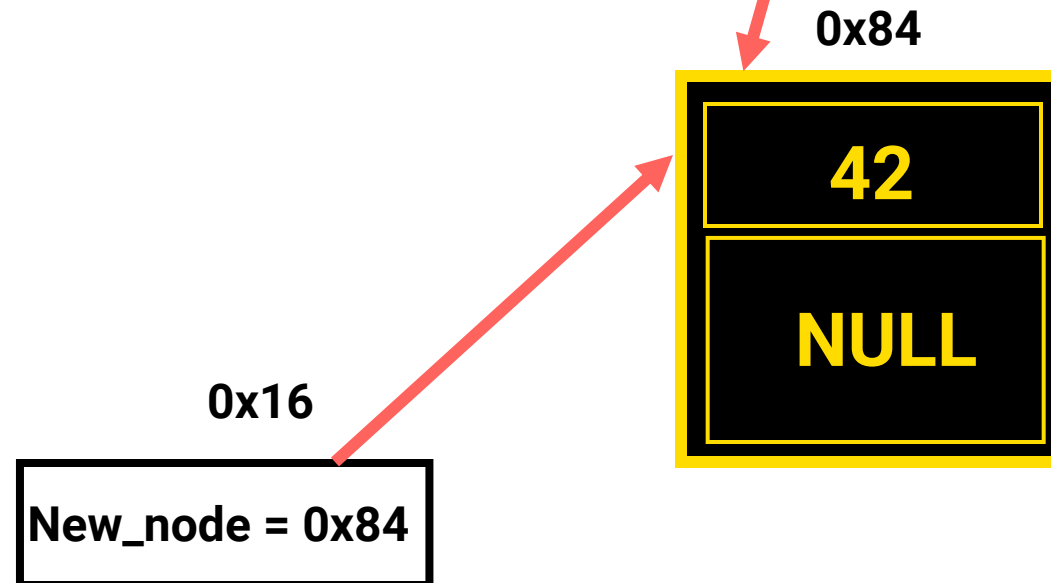


Adding a new node to the end of our list

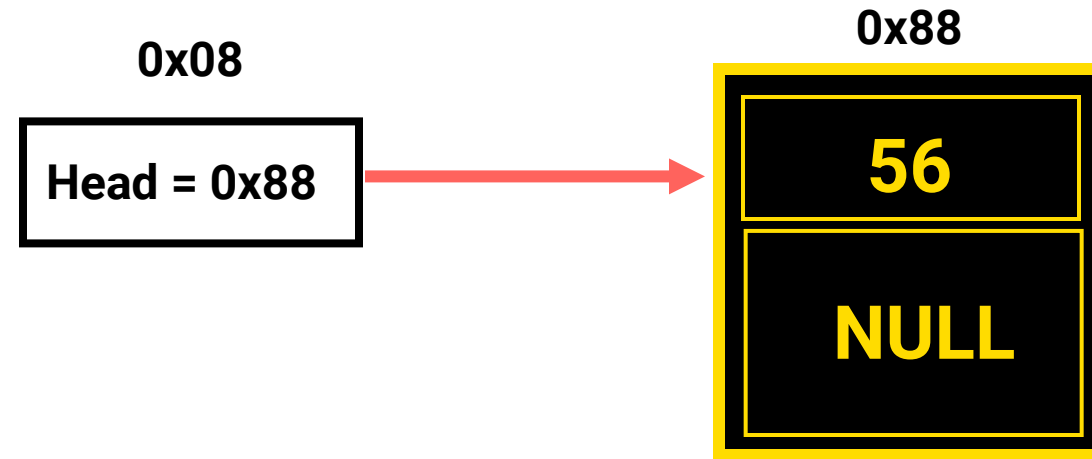
are they connected now? (YES!!)



```
// Connect(link) the head of the list to the new_node  
head->next = new_node;
```



Adding a new node to the start of our list

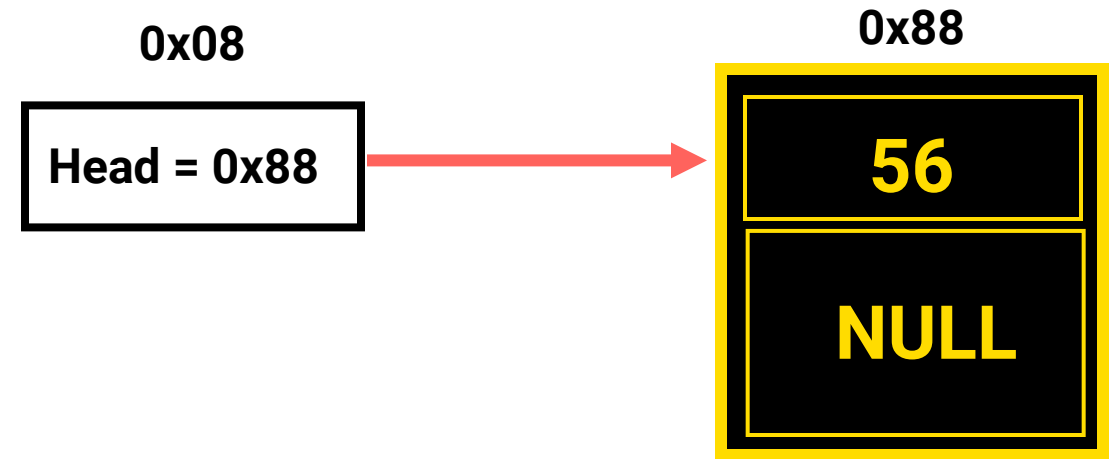


We will create a new node and **insert it at the beginning** of the list.

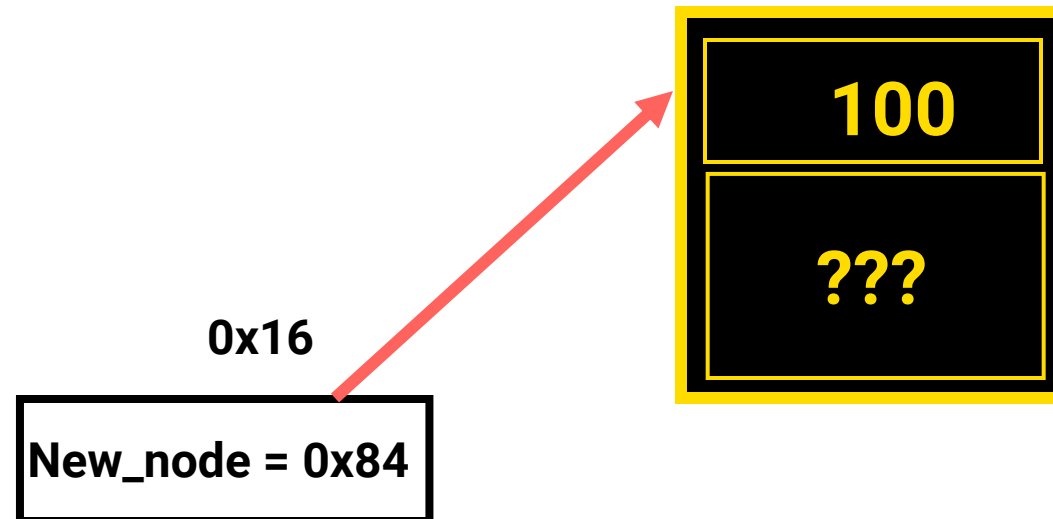
The **first node** in a linked list is commonly referred to as the **head**.

Adding a new node to the start of our list

is the new node connected to the head? (not yet)



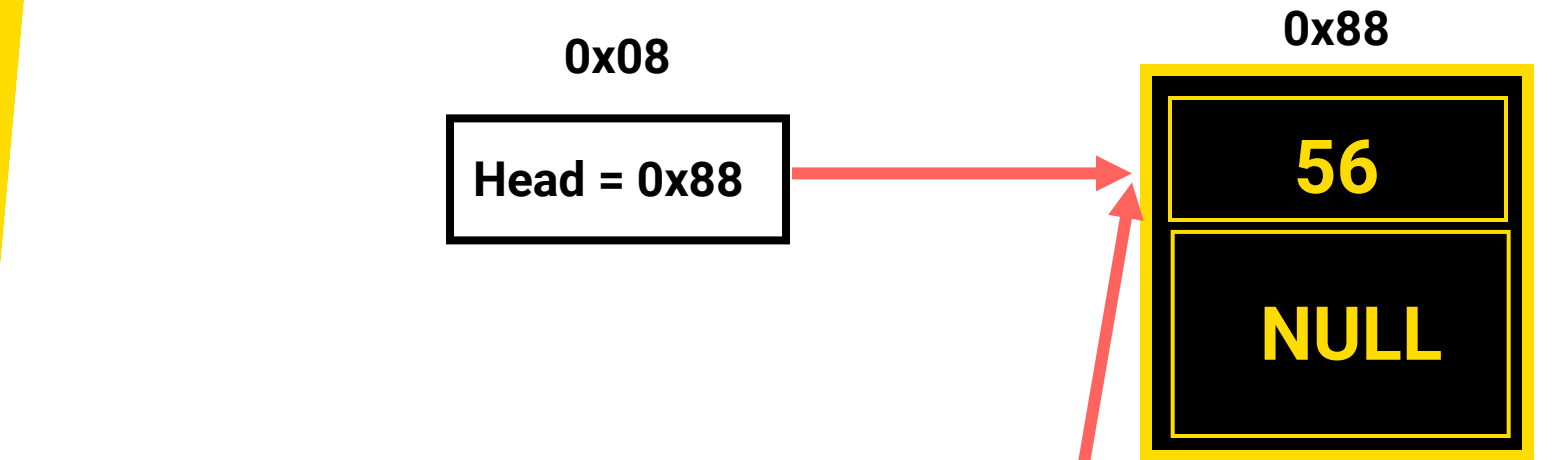
```
struct node *new_node = malloc(sizeof(struct node));  
new_node->data = 100;  
new_node->next = ???;
```



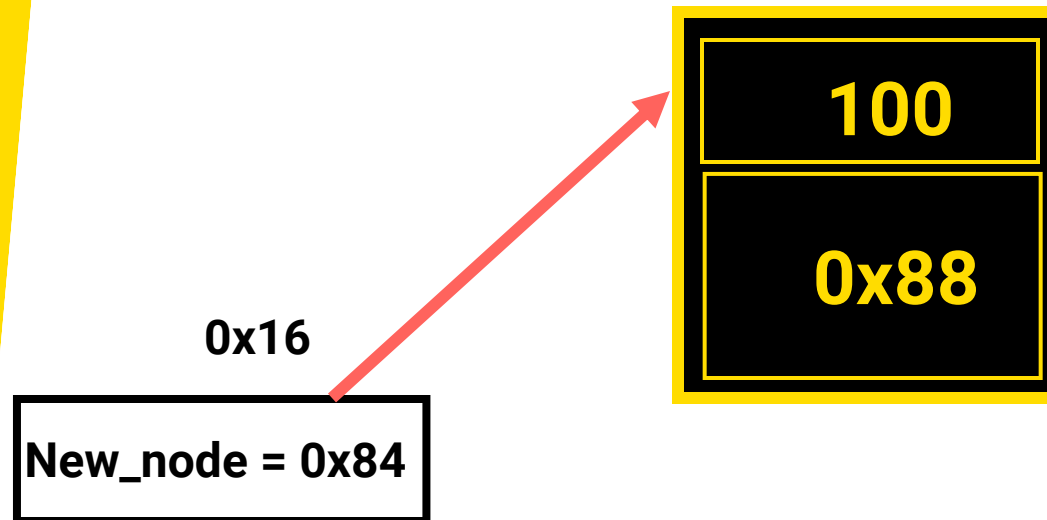
Adding a new node to the start of our list

is the new node connected to the head? **YES!!!**

Is the **head** pointing to the correct node? (not yet)



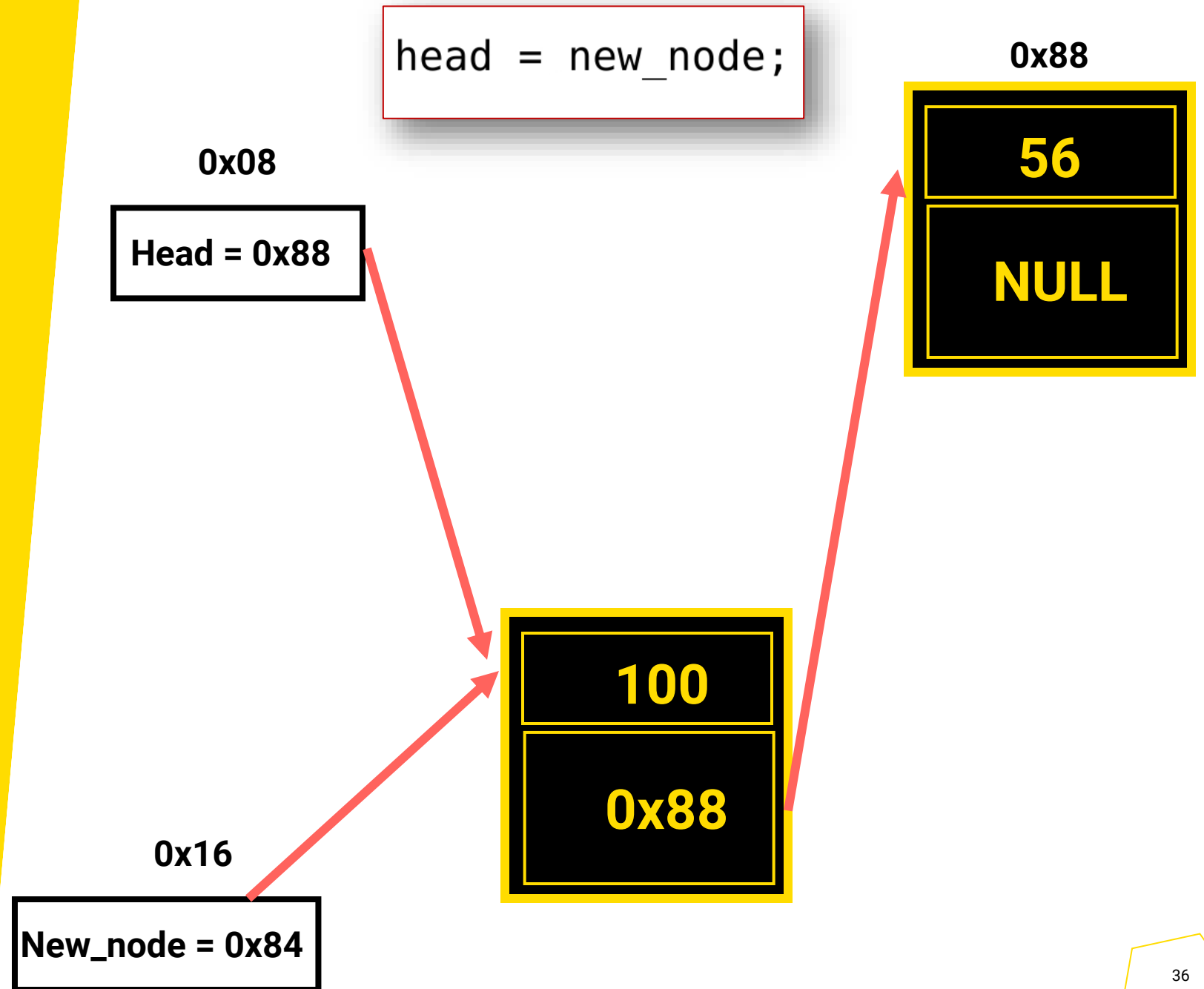
```
struct node *new_node = malloc(sizeof(struct node));  
new_node->data = 100;  
new_node->next = head;
```



Adding a new node to the start of our list

is the new node connected to the head? **YES!!!**

Is the **head** pointing to the correct node? **(YES!!!)**



Linked Lists Functions

How can we move the **insertion** logic into its own **function**?

How could we **repeatedly** call that function to **build a list** from user input?

How would we **search** for a specific value in the list, even if it **contains thousands of nodes**?

How could we **remove** a node from the **middle** of the list safely?

•
•
•
•

We can improve our program by writing functions

e.g. create_node

```
// Allocates memory for a new node and initialises it  
// with the given data value and next pointer.  
// Returns NULL if memory allocation fails.
```

```
struct node *create_node(int data, struct node *next) {  
  
    // Allocate memory for one struct node  
    struct node *new_node = malloc(sizeof(struct node));  
  
    // Check if memory allocation was successful  
    if (new_node == NULL) {  
        return NULL;  
    }  
  
    // Assign the data value to the node  
    new_node->data = data;  
  
    // Set the next pointer  
    new_node->next = next;  
  
    // Return the address of the newly created node  
    return new_node;  
}
```

Create a Linked List Inserting at Head

```
int main(void) {  
  
    struct node *head = NULL;  
    int i = 0;  
  
    while (i < 10) {  
        struct node *new_node = create_node(i, head);  
        head = new_node;  
        i++;  
    }  
  
    return 0;  
}
```

Printing a Node

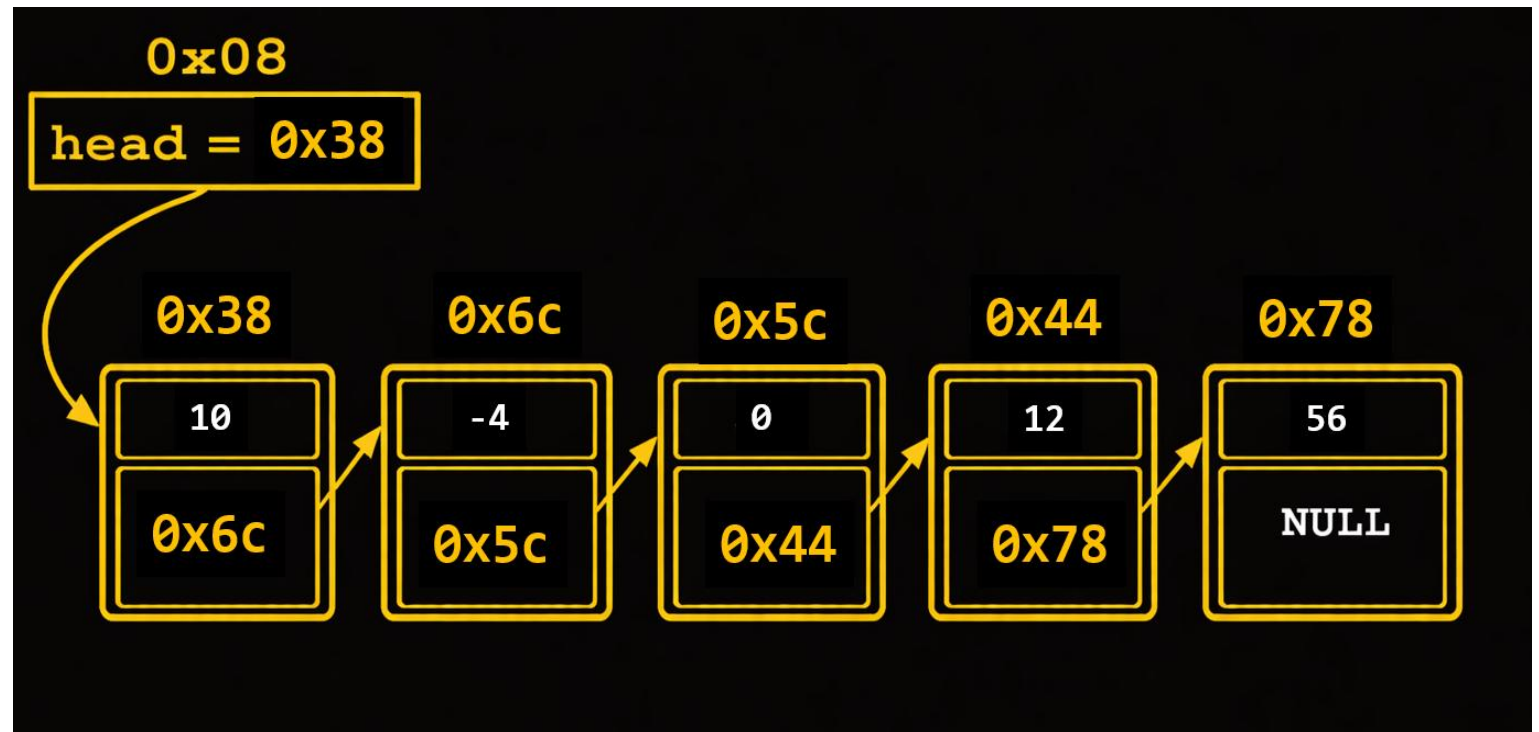
How could we print the data from the **first** node in this linked list?



Printing a Node

How could we print the data from the **first** node in this linked list?

```
printf("%d", head->data);
```



Printing a Node

How could we print the data from **each** node in this linked list?

We need to **traverse** the list



Traversing a Linked List

Traversing a list means **starting at the head node** and **moving through each node one at a time** until reaching the end (**NULL**).

We commonly **traverse** a linked list to:

print the data stored in each node,

determine **how many nodes** are in the list, or

search for a **specific** value within the list.

Traversing a Linked List

First, we should set a **pointer to the beginning of the list**

```
struct node *current = head;
```



Traversing a Linked List

Then we should **move current** to the next nodes one by one



Traversing a Linked List

Then we should **move current** to the next nodes one by one

```
current = current->next;
```



Traversing a Linked List

Then we should **move current** to the next nodes one by one

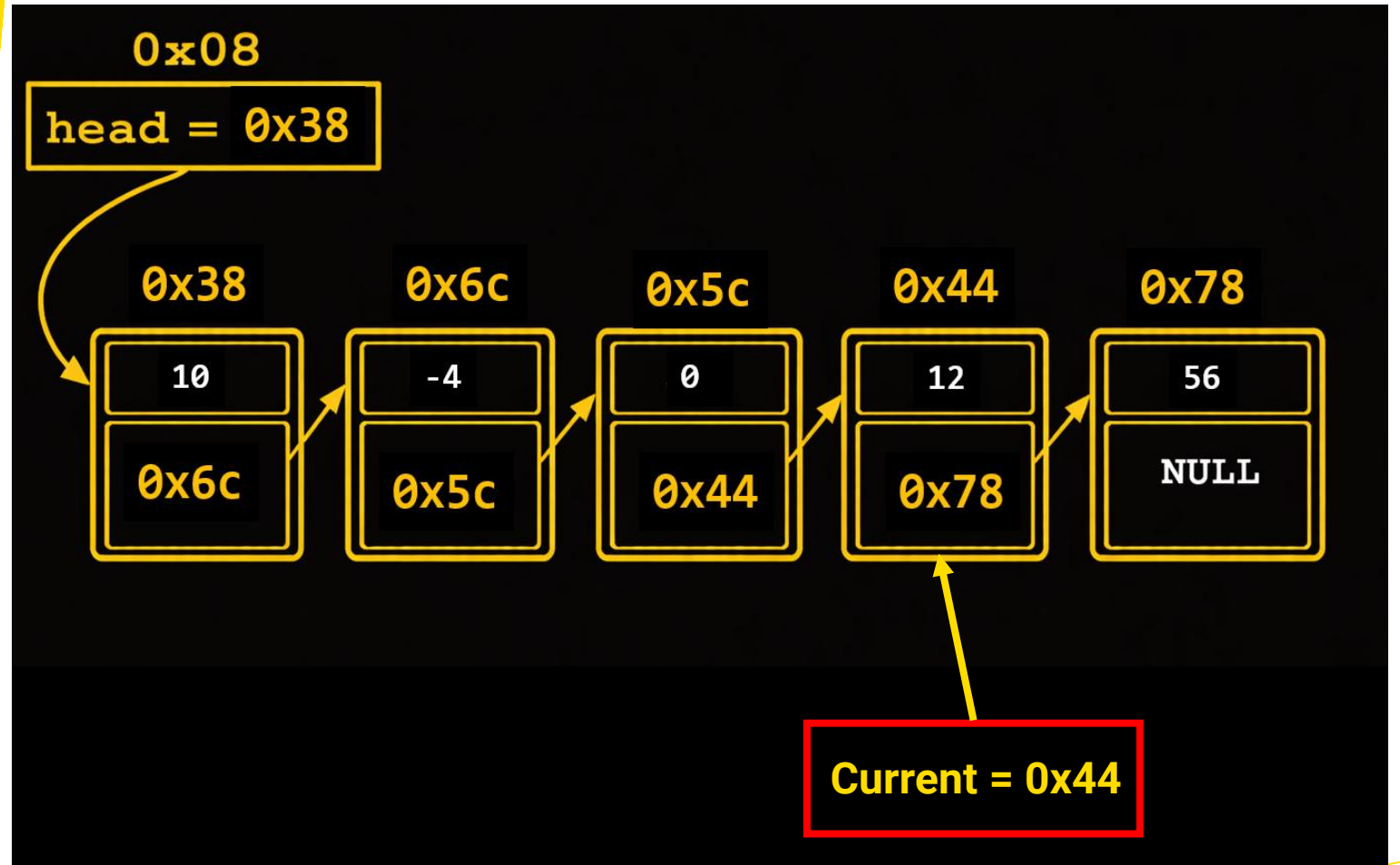
```
current = current->next;
```



Traversing a Linked List

Then we should **move current** to the next nodes one by one

```
current = current->next;
```



Traversing a Linked List

Then we should **move current** to the next nodes one by one

```
current = current->next;
```



Traversing a Linked List

We should stop when
current = NULL

```
current = current->next;
```



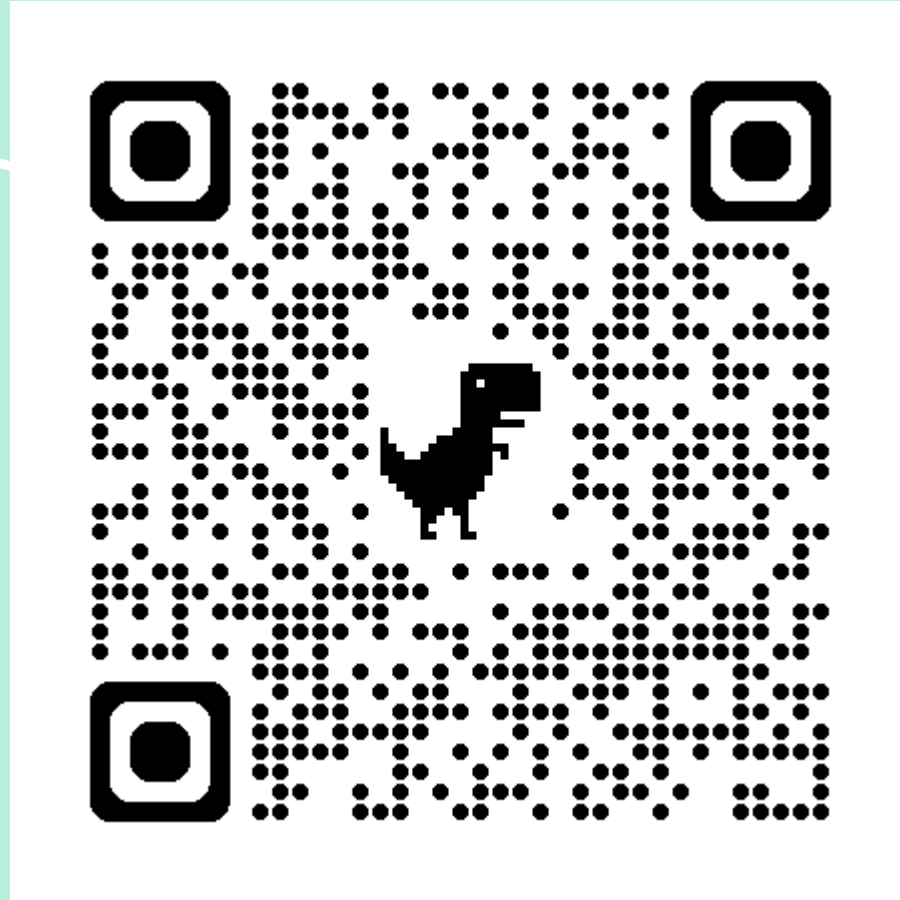
Printing a Linked List

```
// Traverse the list and print the
// data stored in each node
void print_list(struct node *head) {

    struct node *current = head;
    while (current != NULL) {
        printf("%d ", current->data);
        current = current->next;
    }
    printf("\n");
}
```

Demo

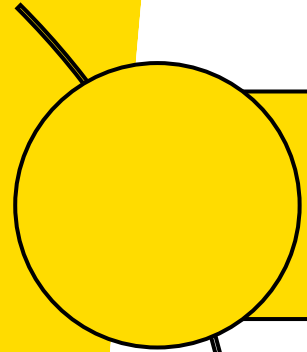
`Print_list_of_three_nodes.c`



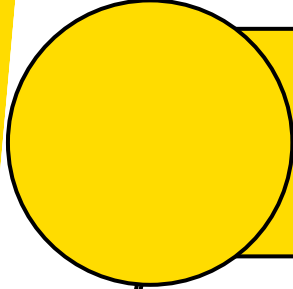
Live lecture code is written for teaching, not perfection.
It may include extra comments and may not always follow
ideal coding style

Inserting nodes in a Linked List

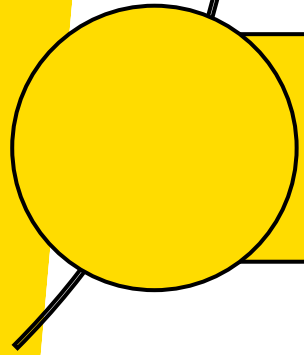
Where can we insert in a list?



At the head (what we just did! Slide 37)



Between any two nodes that exist (next lecture!)



After the tail as the last node (now!)

To add a new node to the **end of a linked list**, we must:

Insert at the tail

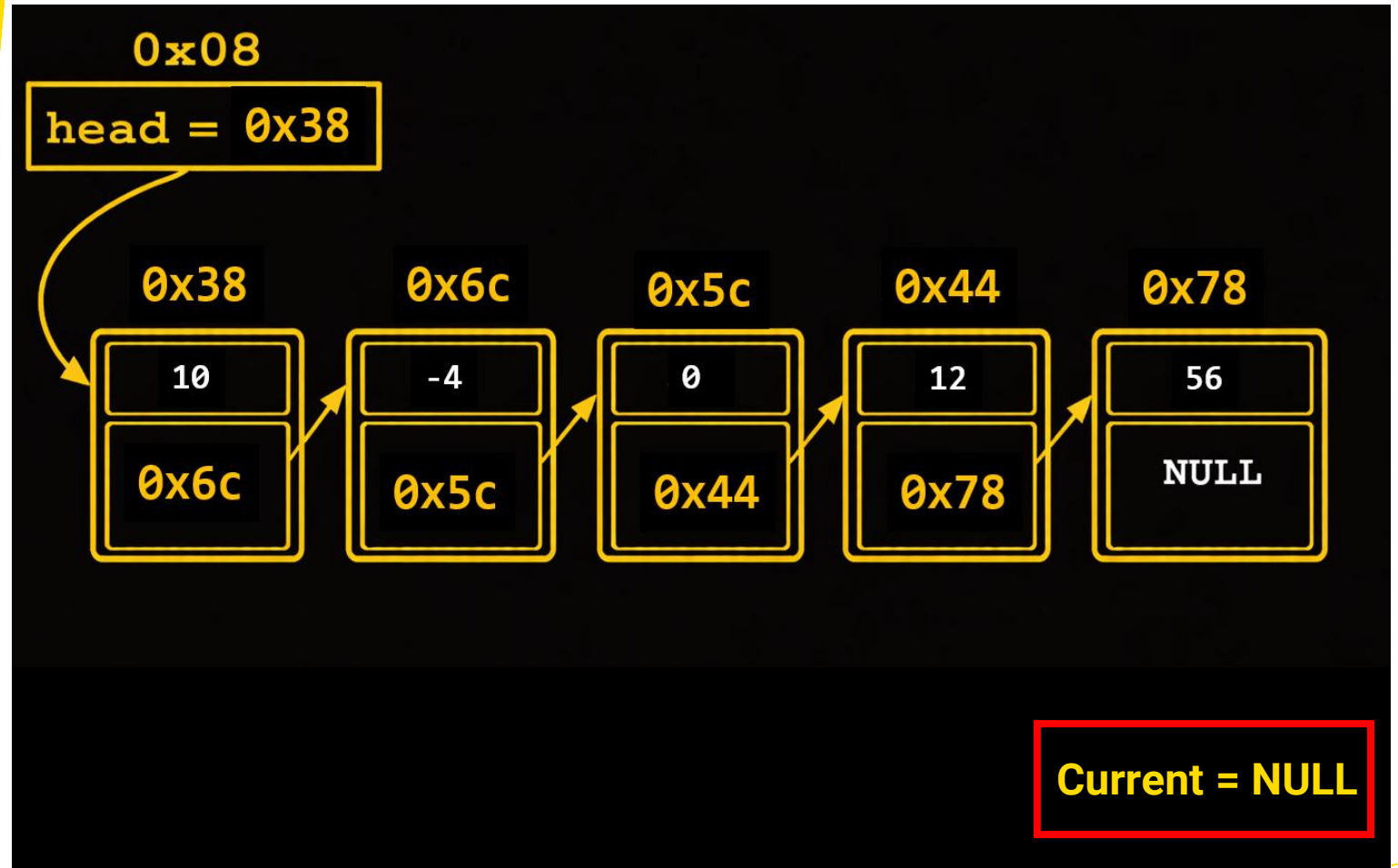
Locate the **last node** in
the list

Update its **next pointer**
so that it **points to the
new node**

How can I move through the list to find the last node?

We should traverse the list

But if we continue **traversing** until `current == NULL`, we have moved **beyond** the **last node** (the tail) of the list!!!



How can I move through the list to find the last node?

We want to stop at the last node

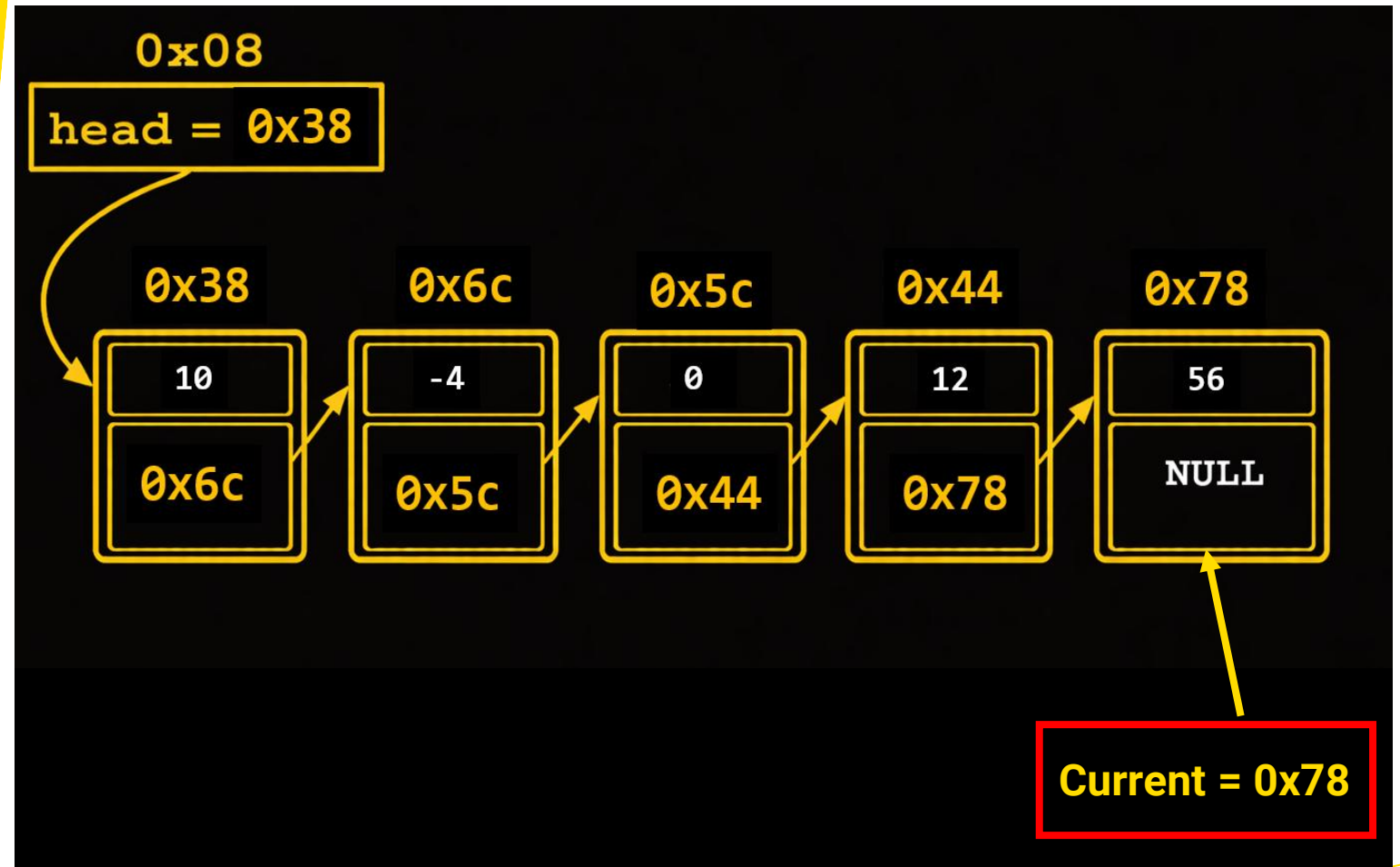
How can we tell if we are at the last node?



How can I move through the list to find the last node?

We want to stop at the last node

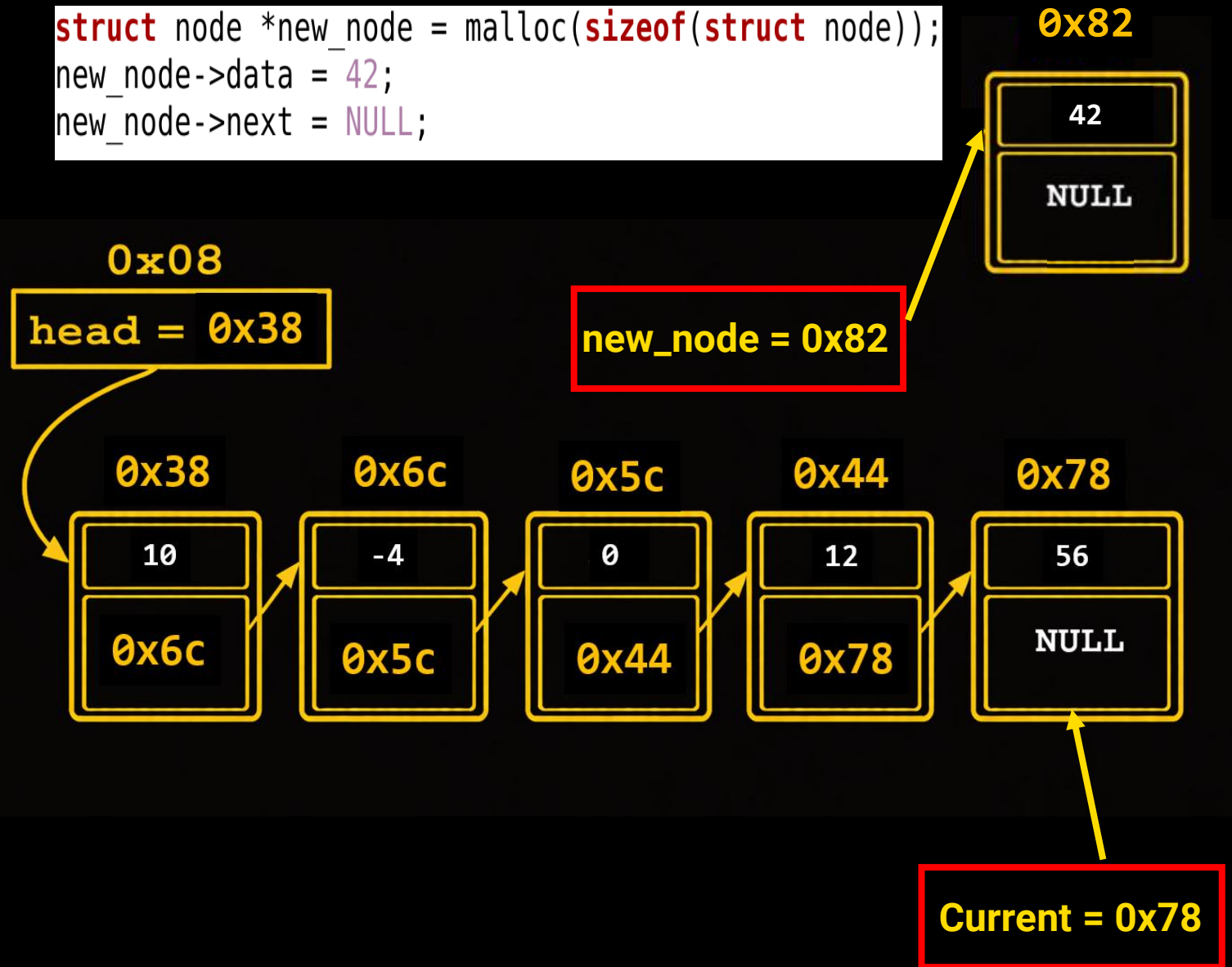
We should stop when `current->next == NULL`



Insert at the tail

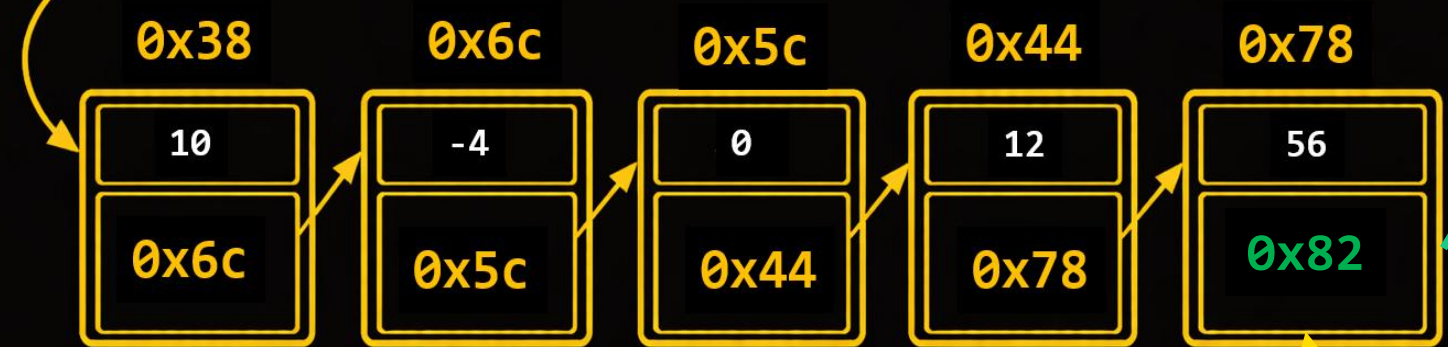
Now we can make the new node first,

```
struct node *new_node = malloc(sizeof(struct node));  
new_node->data = 42;  
new_node->next = NULL;
```

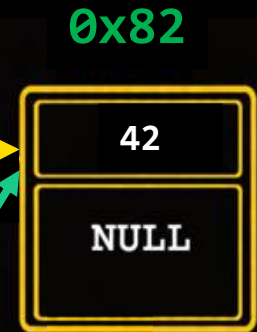


Insert at the tail

0x08
head = 0x38



new_node = 0x82



then link the last node to the new node

current->next = new_node;

Current = 0x78

Insert at the tail function

Big issue

```
// What would cause this function to break?
void insert_at_tail(struct node *head, int data){

    struct node *current = head;

    // Find the tail of the list
    while (current->next != NULL) {
        current = current->next;
    }

    // Create the new node
    struct node *new_node = create_node(data, NULL);

    // link the new node to the tail of the list
    current->next = new_node;

}
```

Testing Linked List Functions

It is essential to **consider different scenarios**:

- An empty list
- A list containing a single node
- A list containing multiple nodes

Our function only inserts nodes **at the end of the list**. If we were developing a function that **inserts a node at any position**, we would **also** need to test:

- Inserting at the beginning of the list
- Inserting in the middle of the list
- Inserting at the end of the list

Inserting At Tail Code Bug

If we have an **empty list** then,

```
head == NULL;  
then current == NULL;
```

it means that **current->next** will be **dereferencing a NULL pointer** and result in a **run time error**

```
void insert_at_tail(struct node *head, int data){  
  
    struct node *current = head;  
  
    // Find the tail of the list  
    while (current->next != NULL) {
```

Inserting At Tail

Is this a better code?

```
void insert_at_tail(struct node *head, int data) {  
    // Create a new node with the given data  
    struct node *new_node = create_node(data, NULL);  
    // Handle the situation where the list is currently empty  
    if (head == NULL) {  
        head = new_node;  
    } else {  
        // Start from the head to traverse the list  
        struct node *current = head;  
  
        // Move through the list until we reach the final node  
        while (current->next != NULL) {  
            current = current->next;  
        }  
  
        // link the new node to the tail of the list  
        current->next = new_node;  
    }  
}
```

Inserting At Tail

The program no longer crashes. However, the list is **still empty** after calling the function. Why does this happen??

```
int main(void) {  
    struct node *head = NULL;  
    insert_at_tail(head, 42);  
    // local variable head is in main is still NULL  
    return 0;  
}
```

How to fix it?

We need to change the function prototype so that it returns the head of the list and then update our local variable by assigning it to the returned value.

```
struct node *insert_at_tail(struct node *head, int data);
int main(void) {

    struct node *head = NULL;

    // updating local variable head
    head = insert_at_tail(head, 42);

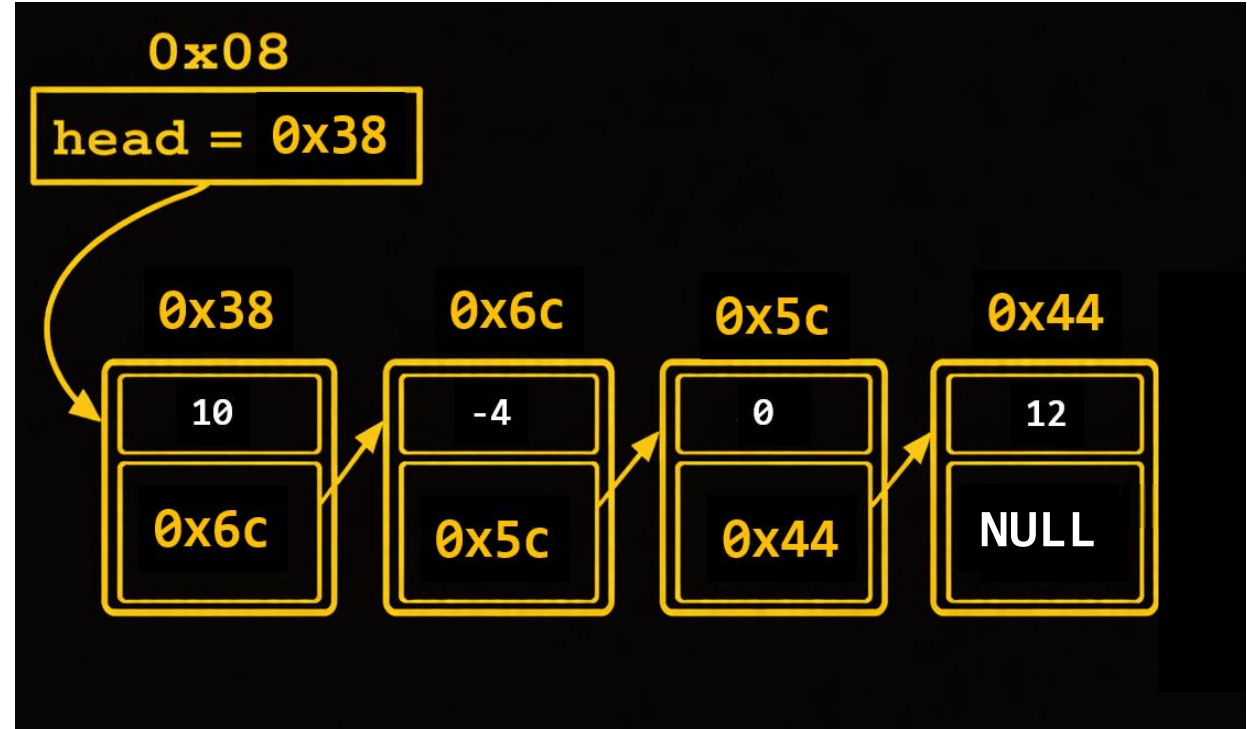
    return 0;
}
```

So, the final Insert at Tail function would be

```
struct node *insert_at_tail(struct node *head, int data) {  
    // Create a new node with the given data  
    struct node *new_node = create_node(data, NULL);  
  
    // Handle the situation where the list is currently empty  
    if (head == NULL) {  
        head = new_node;  
    } else {  
  
        // Start from the head to traverse the list  
        struct node *current = head;  
  
        // Move through the list until we reach the final node  
        while (current->next != NULL) {  
            current = current->next;  
        }  
  
        // link the new node to the tail of the list  
        current->next = new_node;  
    }  
    return head;  
}
```

Inserting in the Middle of the List

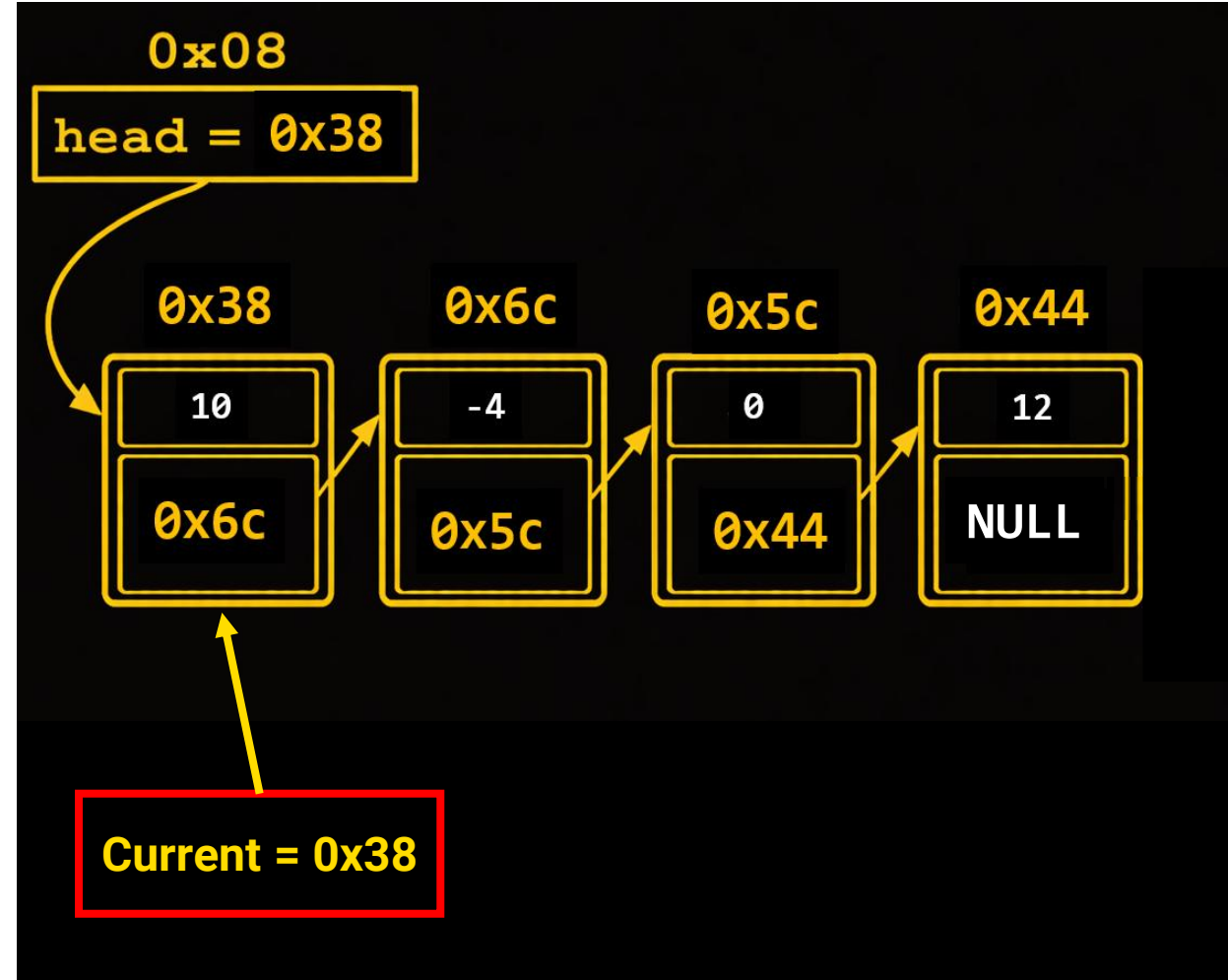
We want to **insert a new node** at the **position given by size / 2**, assuming positions **begin counting from 0**. In this example, that means the new node should be inserted at position 2.



This means, we should **use a counter** and **stop traversing** when we get to the node **before the position** we want to insert at (**size/2 - 1**). In this case position 1.

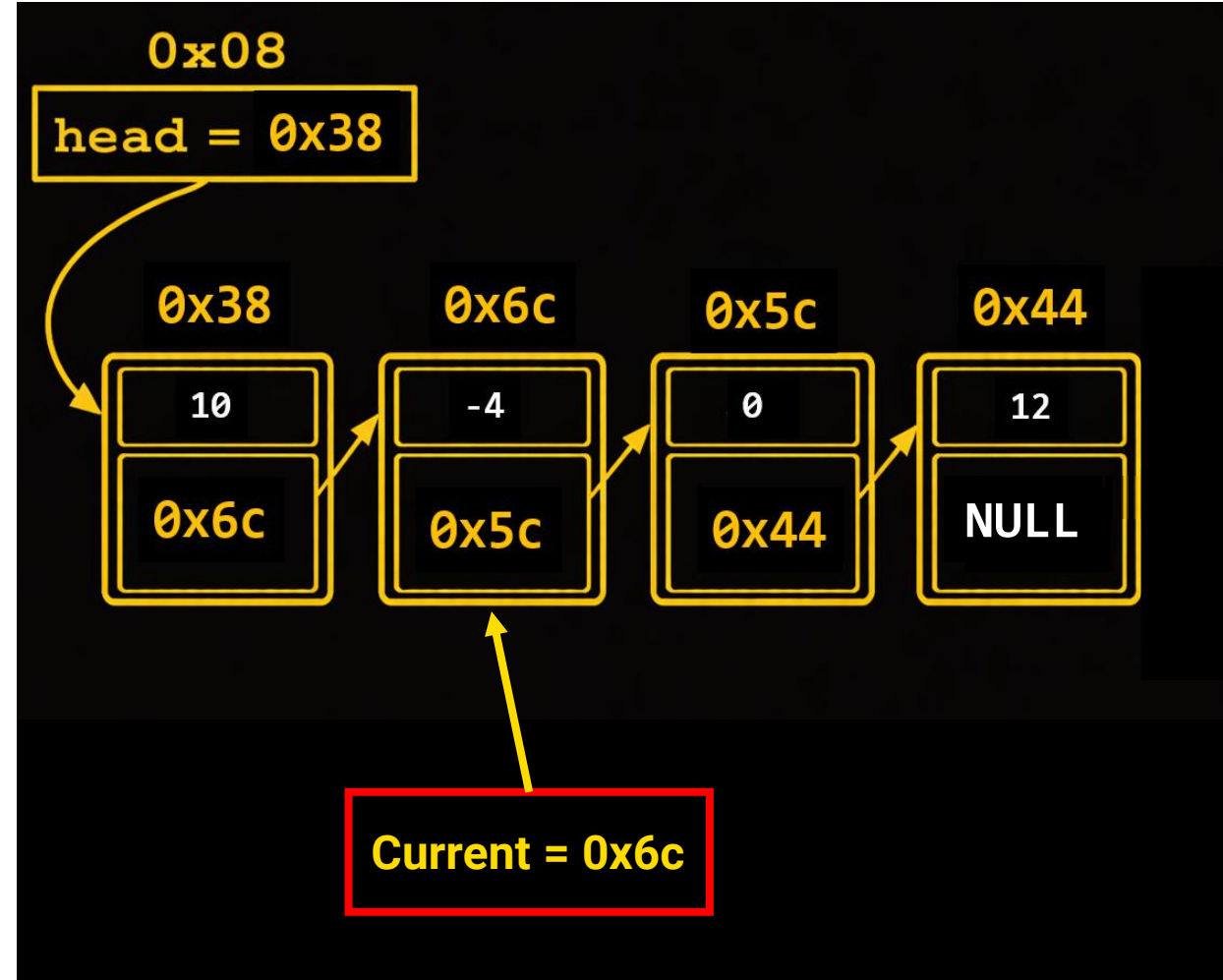
Inserting in the Middle of the List

```
struct node *current = head;  
int counter = 0;
```



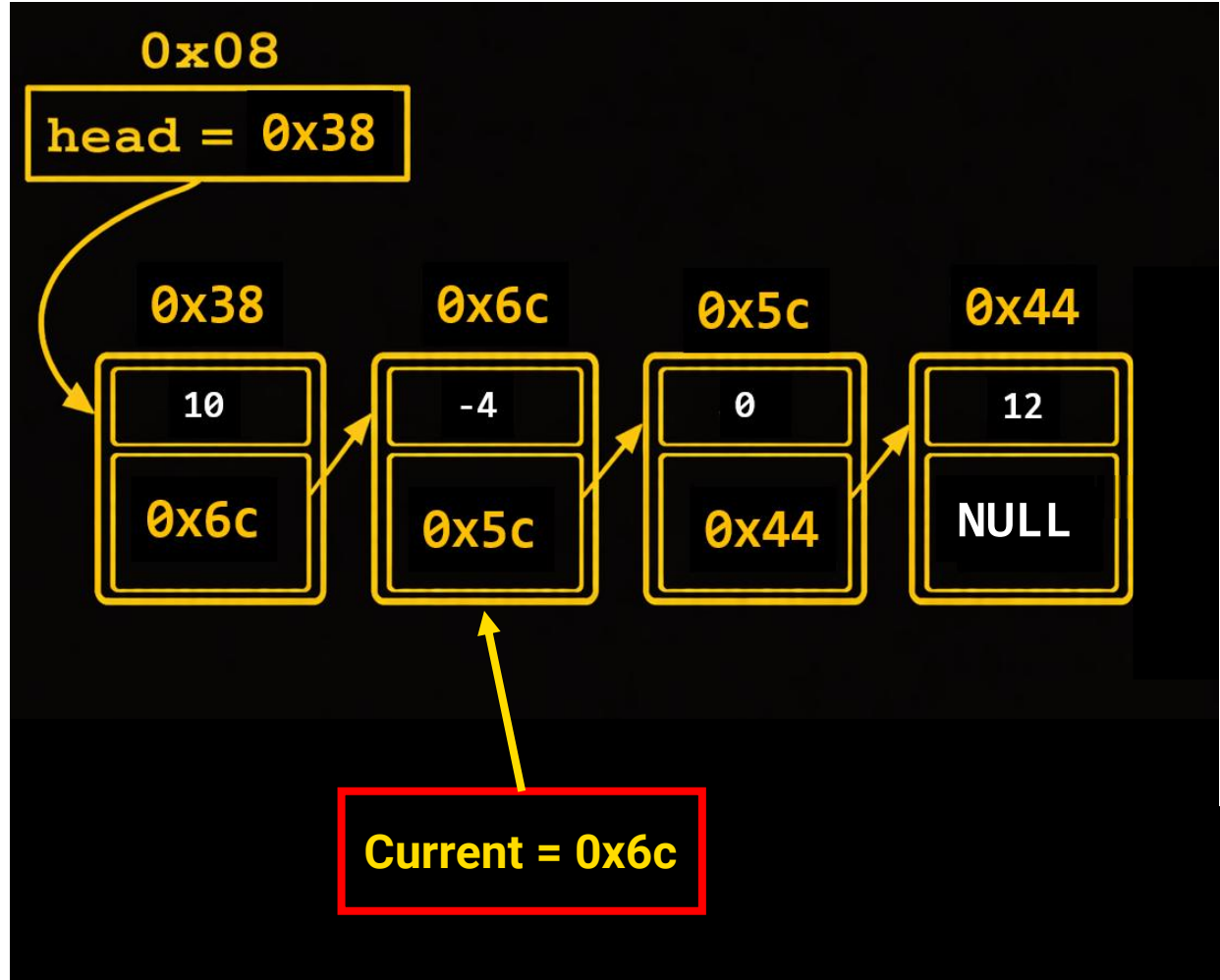
Inserting in the Middle of the List

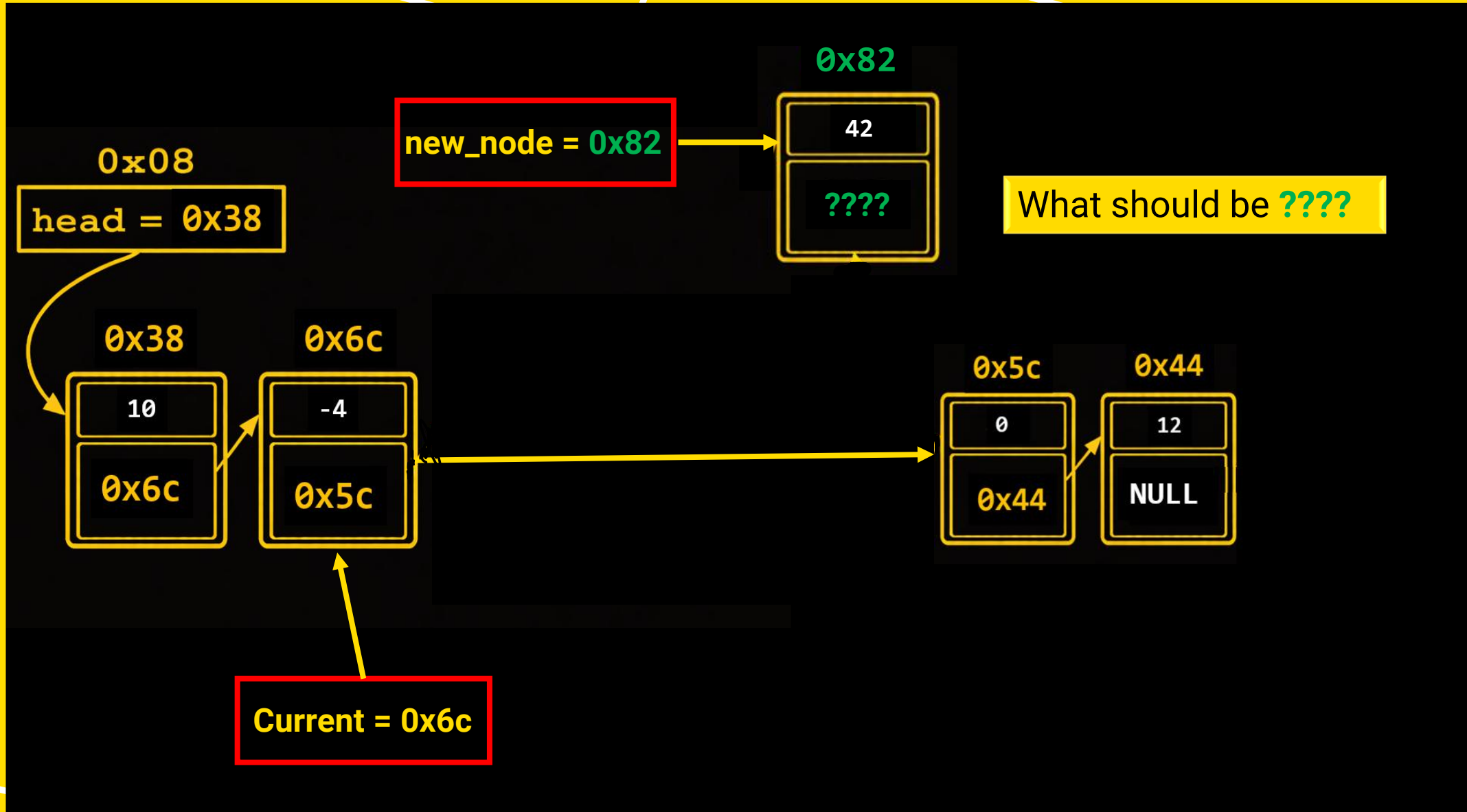
```
while (counter < size/2 - 1) {  
    current = current->next;  
    counter++;  
}
```

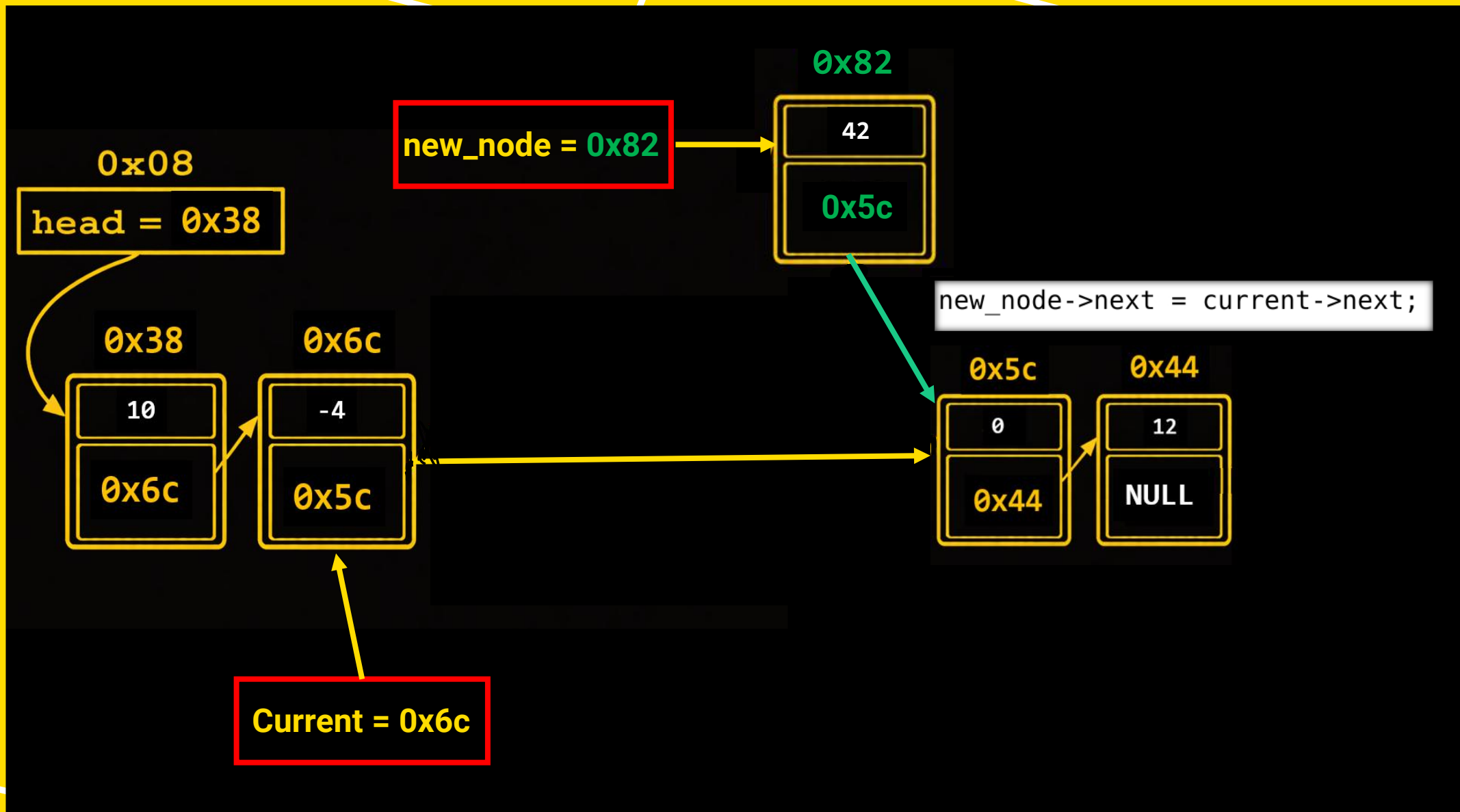


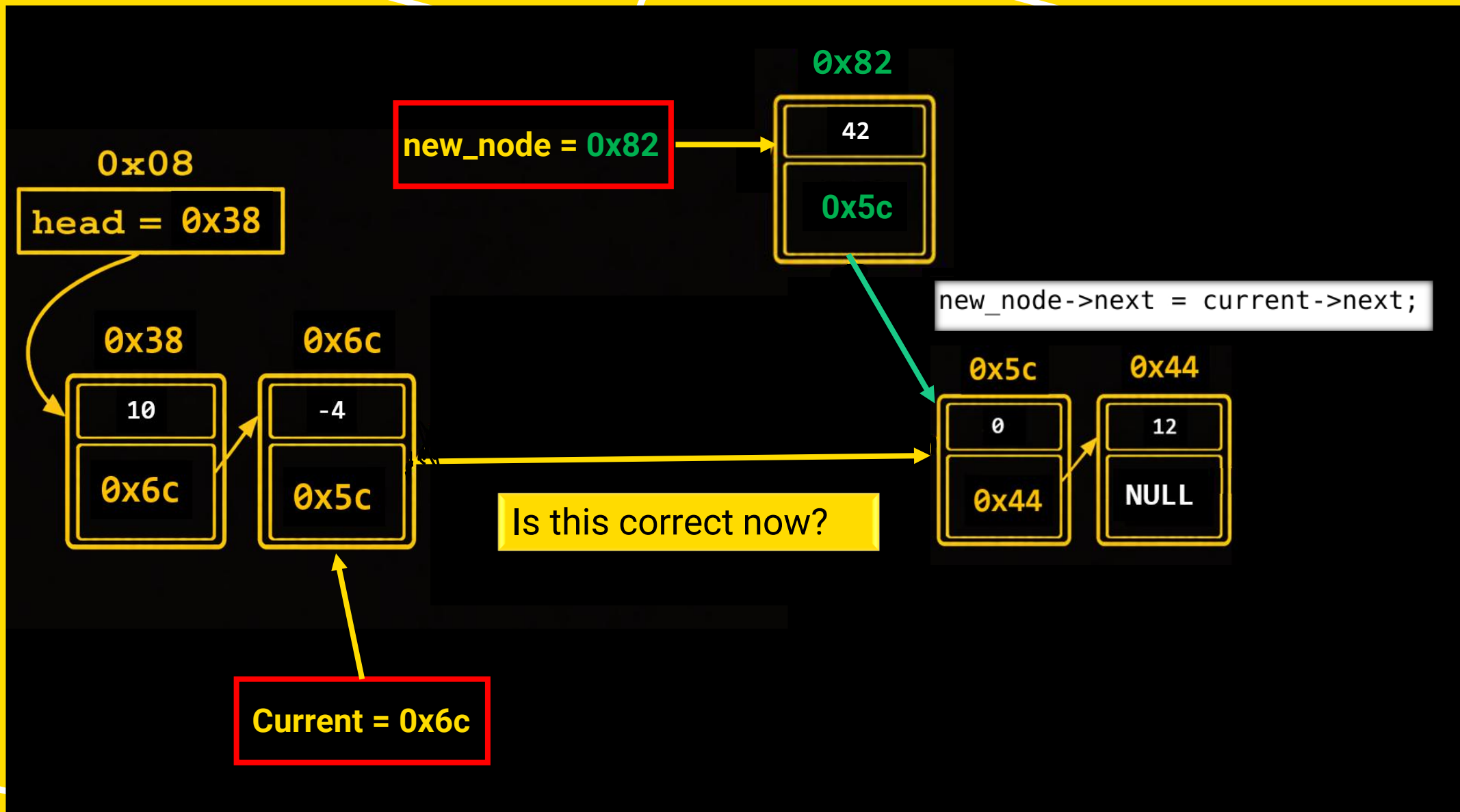
Now we want to connect our **new node**. It should come **after the current node**, but **before current ->next**

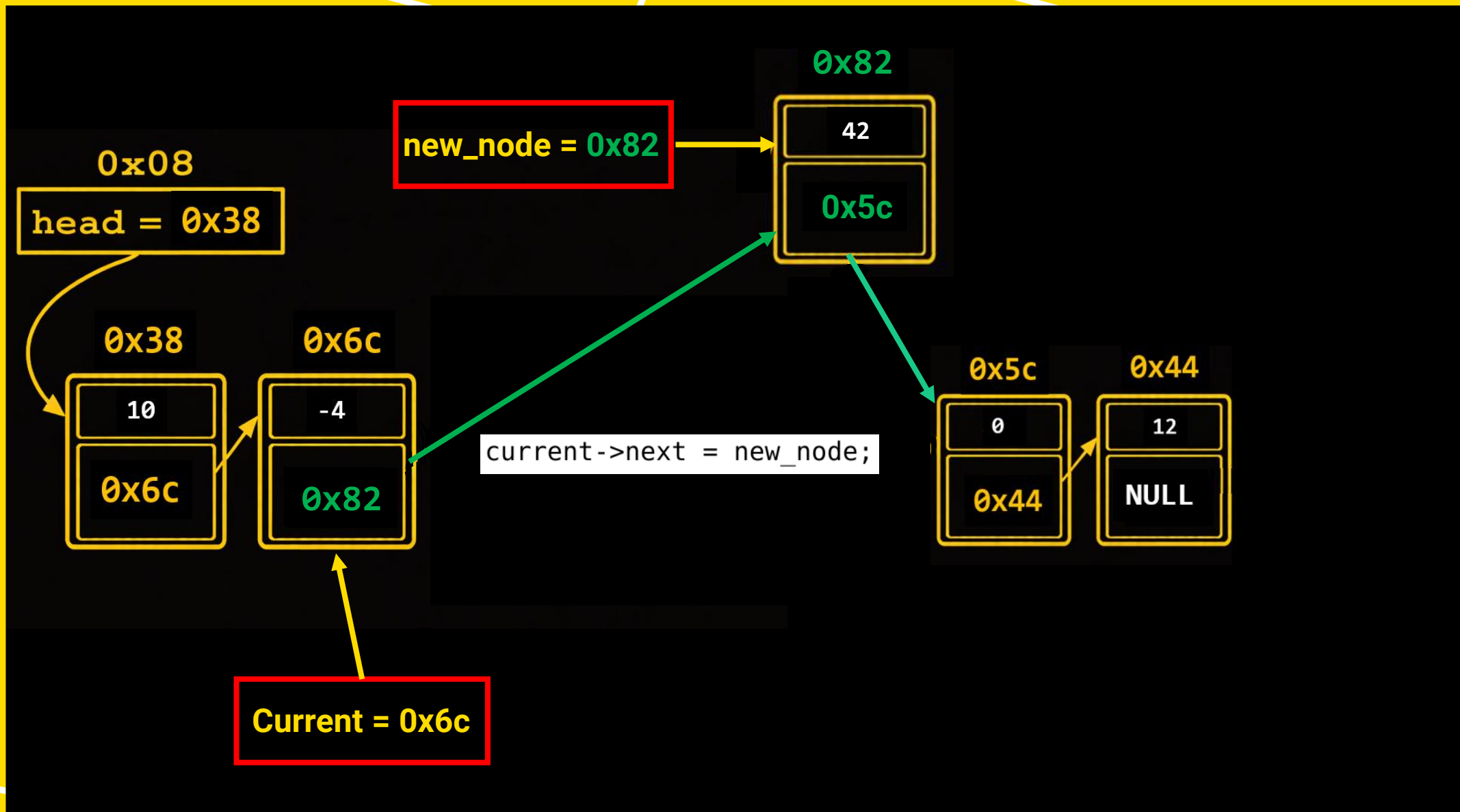
Inserting in the Middle of the List











What would be the coding issues?

What conditions will break this?

What happens if it is an **empty list**?

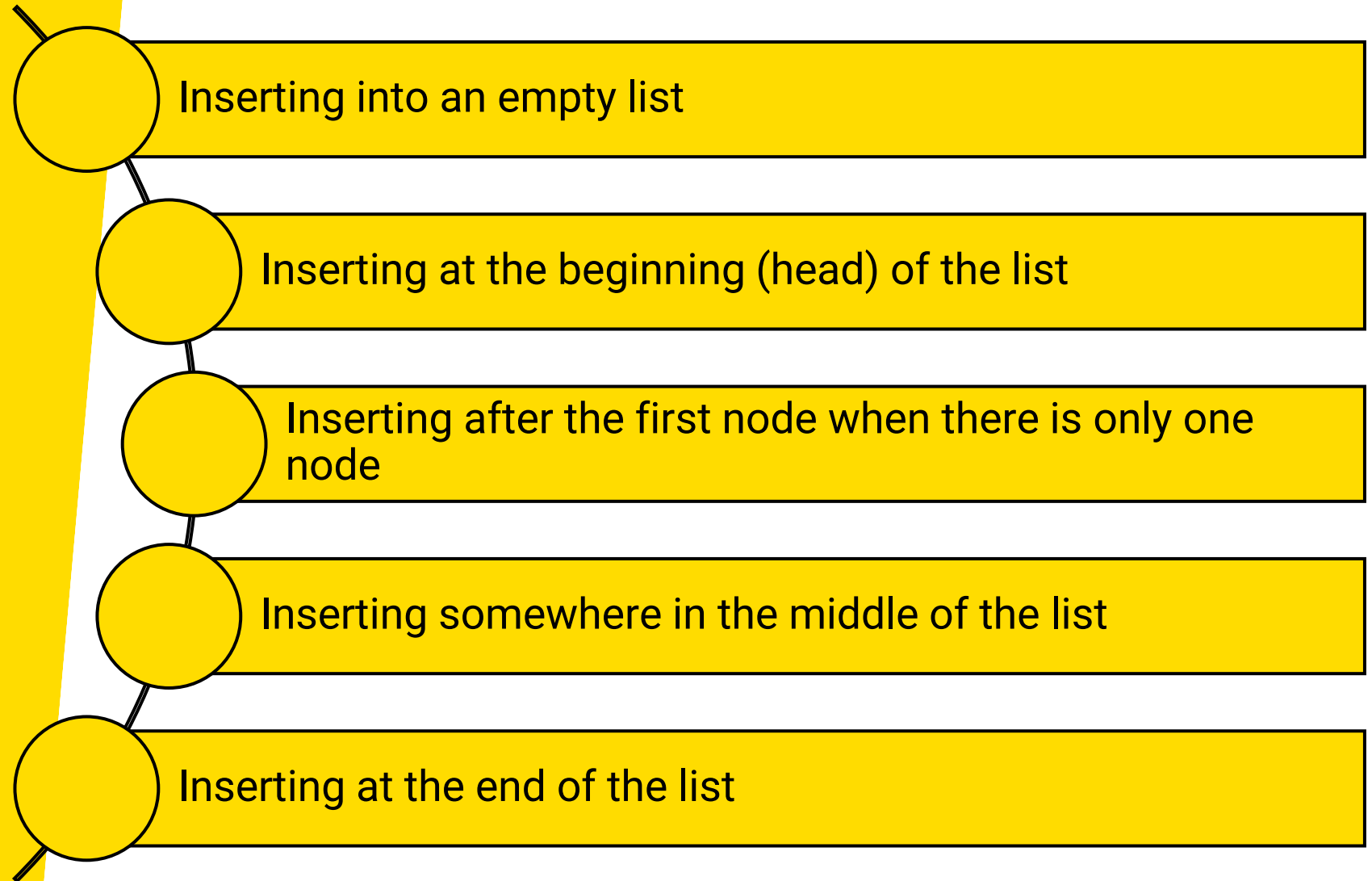
What happens if there is **only 1 item** in the list?

Anything else we should check?

How can we modify our code to handle any of these situations that break it?

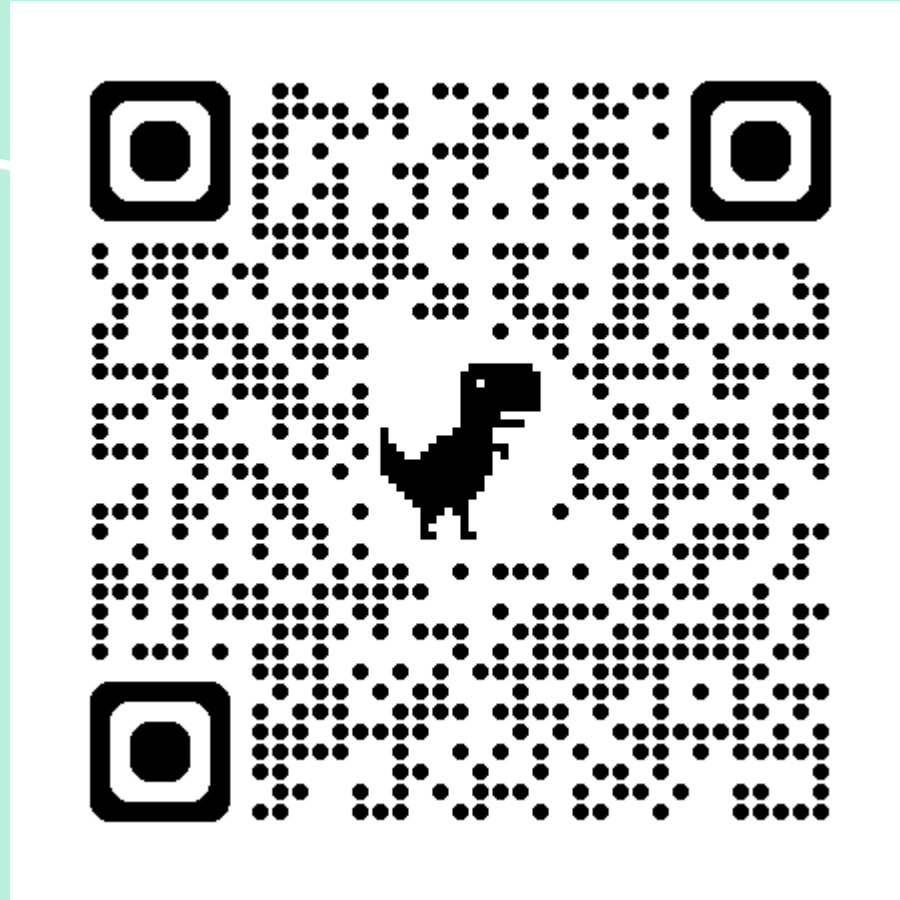
Inserting Into a Linked List Test Cases

Draw a
diagram!!!



Demo

`list_create_insert_print_length.c`



Live lecture code is written for teaching, not perfection.
It may include extra comments and may not always follow
ideal coding style

Voice of the Student

Anonymous ongoing feedback
Anything you wanted to share with me



26T1 Voice of the Student



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See you soon ...